# Math 312: Introduction to Number Theory Lecture Notes

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These are rough notes for the summer 2011 course. Problem sets an internal website.	and solutions were posted on

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#### Introduction

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### 0.1. Cold open

• Theory of "numbers", mainly meaning whole numbers, that is the integers.

#### Cryptography.

- Do you use the internet?
- Can use number theory to establish identity ("The person who knows the factorization N = pq")
  - Key (1): There are arithmetic problems that only the person who knows the factorization can solve.
  - Key (2): I can prove to you that I know to factor N without revealing this number.
- How do you know that https://www.yourbank.ca is really controlled by your bank? Roughly speaking:
  - The manufacturer of your browser equips it with a list of 230 numbers and told to trust those people who know how to factor them. Those people are called "certificate authorities". For simplicity assume there was only one CA, with its number  $N_{\text{CA}}$  built into the browser.
  - Your bank creates a number  $N_{\rm bank}$  for itself. It goes to the CA and gets a "digital certificate", which says: "the people who know how to factor  $N_{\rm CA}$  say that the server at https://www.yourbank.ca knows how to factor the number  $N_{\rm bank}$ ". Moreover, the certificate includes a number calculated using the details in it and also  $p_{\rm CA}, q_{\rm CA}$  so it cannot be forged.
  - When you access https://www.yourbank.ca, the website provides your browser with the certificate. Your computer can *verify* the signature using the number  $N_{\rm CA}$  that it knows. If this is ok it then challenges the machine on the other side to prove that it can really factor the number  $N_{\rm bank}$  specified in the certificate.
  - If that works too the browser is happy.
  - (I'm Lying a little here) When you type your password on https://www.yourbank.ca, what your browser sends is not simply the characters you type (that would be bad if there were eavesdroppers). What it sends instead is the result of a calculation involving the number  $N_{\text{bank}}$ . The calculation is done in such a way that the bank can check that you typed the correct password using its secret knowledge ( $p_{\text{bank}}$ ,  $q_{\text{bank}}$ ) whereas the eavesdropper (who only knows  $N_{\text{bank}}$ ) can't learn what the password was.

There are practical problems with this scheme, but we won't discuss internet security. What we will discuss is the *number theory* that makes secure websites possible.

Summary for those reading the notes:

- Underlying fact of life: arithmetic is "easy" but solving equations is "hard".
- There is a method for me to convince others that I know p,q without revealing them. This allows me to prove my identity ("I'm the person who knows to factor N"). But prevents forgery
- The method extends to verifiable signatures ("This email was written by the person who knows how to factor N"). No-one can forge my signature, but they can check I'm the person who signed (signing requires knowing p,q but checking only requires knowing N).
- The method extends to secure communication: you can take a message m and send me F(m,N). Knowing p,q solving F(m,N) = C is actually easy. But it is believed impossible to solve the equation knowing only N.
- We will discuss how to do all these things (signature protocol, secure communications protocol).

# Other examples

• Software "product activation" key verification.

#### 0.2. Administrivia

Syllabus distributed. Other key points:

- Problem sets will be posted on the course website. Solutions will be posted on a secure system (email explanation will be sent).
  - Depending on time, the grader may only mark selected problems. Solutions will be complete.
- Absolutely essential to
  - Read ahead according to the posted schedule. Lectures after the first will assume that you had done your reading.
  - Do homework.
- Office hours on Tuesdays 2-3 (not on May 10th) and Thursdays 11-12. Also by appointment
- Course website is important. Contains notes, problem sets, announcements, reading assignments etc.

#### **0.3.** More on problems of number theory

#### Some arithmetic.

- It is now 10am. What hour will it be 5 hours from now?
- Today is Tuesday, May 10th. What day was it 4 days ago?
  - What day of the week did 10/5/2010 fall on? Hint:  $365 = 52 \cdot 7 + 1$ .
  - What about July 1st, 1867? [don't wait for answer]

#### Properties of individual numbers.

• Can we find integers a, b so that  $\left(\frac{a}{b}\right)^2 = 2$ ?

- We want to at least find integers a, b so that  $\frac{a}{b}$  is close to  $\sqrt{2}$ . How well can we do this without taking b too large?
- Consider the decimal expansion of  $\pi$  (or e, or your favourite number). Do all the digits  $0, 1, \dots, 9$  appear  $\frac{1}{10}$ th of the time? What about all sequences  $00, 01, \dots, 99$ ?
- We already saw that primes are useful. We need to make primes.
  - How many primes are there? Are they easy to find?
  - Specifically: how does one tell if a given number n is prime?

### Main focus: finding integer solutions to equations.

- All integer solutions to  $x^2 + y^2 = z^2$  known to the Greeks (lots!)
  - Only obvious integer solutions to  $x^4 + y^4 = z^4$  (Fermat)
  - Only obvious solutions to  $x^3 + y^3 = z^3$  (Euler)
  - **–** ...
  - Only obvious solutions to  $x^n + y^n = z^n$ ,  $n \ge 3$  (Ribet, Wiles)
- Other equations to solve:
  - -xy = N ("factorization"). Very important we will discuss this.
  - -y = x + d, z = x + 2d, w = x + 3d ("arithmetic progression")
    - \* Szemeredi: if  $A \subset \{1, ..., N\}$  is large enough then for some  $d \neq 0$ , A it contains x, y, z, w solving the above equation.
    - \* Green, Green-Tao: can take A to be the set of primes, or even a "dense subset" of the primes.
  - $p_1 p_2 = 2$  ("twin primes")
  - $p_1 + p_2 = 2N$  ("Goldbach Conjecture")
- Counting solutions to equations
  - # $\{1 \le p \le x \mid p \text{ prime}\} \approx \frac{x}{\log x}$  (Hadamard, de la Vallée-Poussin 1896 following Riemann 1859)
  - $\left| \# \{ 1 \le p \le x \mid p \text{ prime} \} \frac{x}{\log x} \right| \le C\sqrt{x} (\log x)^{100}$  ("Riemann hypothesis")

#### Much more.

Not for today.

#### **0.4.** Course plan (subject to revision)

- The integers
  - Definition: induction and the well-ordering principle;
  - Multiplication: divisibility and the GCD, primes and unique factorization.
- Congruences [major point of the course]
- Algebra: primitive roots
  - The multiplicative group
  - Primality testing
  - Discrete log cryptosystems
- Application: public-key cryptography, RSA
- Multiplicative functions
- Quadratic reciprocity

References. Any book with the title "Elementary Number Theory" or "Introduction to Number Theory" will cover the material. I will generally follow the textbook "Elementary Number Theory and its applications" by K. Rosen.

#### CHAPTER 1

# The Integers

### 1.1. The axioms; the well-ordering principle; induction

We record here the usual properties of the integers.

DEFINITION 1. The *integers* are the sextuple  $(\mathbb{Z}, +, \cdot, 0, 1, <)$  satisfying:

- $+: \mathbb{Z} \times \mathbb{Z} \to \mathbb{Z}$  ("addition") and  $\cdot: \mathbb{Z} \to \mathbb{Z}$  ("multiplication") are binary operations;  $0, 1 \in \mathbb{Z}$  are elements ("zero", "one") and < is a binary relation ("less than").
- Addition is commutative and associative; a + 0 = a and for all  $a \in \mathbb{Z}$  there is  $(-a) \in \mathbb{Z}$  so that a + (-a) = 0.
- Multiplication is commutative and associative;  $a \cdot 1 = a$ .
- The relation < is transitive, and for all  $a, b \in \mathbb{Z}$  exactly one of a < b, a = b, b > a holds.
- For a, b > 0 we have  $a + b, a \cdot b > 0$ .
- Well-ordering property: every non-empty subset  $A \subset \mathbb{N} = \mathbb{Z}_{>0}$  has a least element.

Note that we interchangeably write  $\mathbb{N}$  or  $\mathbb{Z}_{\geq 0}$  for the set of *natural numbers*, those being the non-negative integers (0 is a natural number for us). The textbook also uses the notation  $\mathbb{Z}^+$  for the set  $\mathbb{Z}_{\geq 0} = \mathbb{Z}_{\geq 1}$  of *positive* integers.

The most important property is the last one. We illustrate it with several calculations. We first justify the last equality:

LEMMA 2. (Discreteness) There is no integer b satisfying 0 < b < 1.

PROOF. Let  $A = \{n \in \mathbb{Z} \mid 0 < n < 1\} \subset \mathbb{Z}_{\geq 0}$ . Assume by contradiction that A is non-empty and let a be its minimal element. Then 0 < a < 1. Multiplying both sides by a we find that  $0 < a^2 < a < 1$ . But then  $a^2$  is an integer satisfying  $0 < a^2 < 1$  so  $a^2 \in A$  and  $a^2 < a$ , a contradiction to a being the smallest member of A.

COROLLARY 3. For any integer n there is no integer a satisfying n < a < n + 1.

PROOF. If a existed set b = a - n. Then 0 < b < 1.

THEOREM 4. (Principles of induction)

- (1) (Weak induction) Let  $P \subset \mathbb{N}$ . Assume that  $0 \in P$  and that for any number  $n, n \in P \Rightarrow n+1 \in P$ . Then  $P = \mathbb{N}$ .
- (2) (Strong induction) Let  $P \subset \mathbb{N}$ . Assume that for any  $n \in \mathbb{N}$ ,  $\{a \in \mathbb{N} \mid a < n\} \subset P \Rightarrow n \in P$ . Then  $P = \mathbb{N}$ .

PROOF. (1) Assume by contradiction that  $P \neq \mathbb{N}$ . Then the set  $A = \mathbb{N} \setminus P$  of counterexamples is non-empty. Let m be its minimal element. Then  $m \neq 0$   $(0 \in P)$  so  $m \geq 1$  and  $m-1 \in \mathbb{N}$ . Since m was minimal,  $m-1 \notin A$  so  $m-1 \in P$ . But then  $m = (m-1)+1 \in P$ , a contradiction.

(2) Let Q be the set of  $n \in \mathbb{N}$  such that  $\{a \in \mathbb{N} \mid a < n\} \subset P$ . Then  $0 \in Q$  (the set of natural numbers smaller than zero is empty). Also, if  $n \in \mathbb{N}$  then  $\{a \in \mathbb{N} \mid a < n\} \subset P$  so  $n \in P$ . But then

 $\{a \in \mathbb{N} \mid a < (n+1)\} = \{a \in \mathbb{N} \mid a < n\} \cup \{n\} \subset P \text{ so } n+1 \in Q.$  By part (1) it follows that  $Q = \mathbb{N}$ . Now for any  $n \in \mathbb{N}$  we have that  $n+1 \in Q$  and so that  $n \in P$ .

DEFINITION 5. Call an integer n even if it is of the form n = 2k for some  $k \in \mathbb{Z}$ . Call it odd otherwise.

THEOREM 6. (Division by 2) Among every two consecutive natural numbers at least one is even.

PROOF USING WELL-ORDERING. Let A be the set of  $n \ge 0$  such that n, n+1 are both odd, and let m be a minimal member of A. Then both m, m+1 are odd so m>0 (zero is even!) and  $m-1 \ge 0$ . Since m is minimal,  $m-1 \notin A$ , one of m-1, m is even. m is odd so m-1=2k for some  $k \in \mathbb{Z}$ . But then m+1=(m-1)+2=2(k+1) is even, a contradiction. Since A cannot have a least member it is empty.

PROOF USING WEAK INDUCTION. Let P be the set of  $n \ge 0$  such that at least one of n, n+1 is even.  $0 \in P$  since it is even. Assume that  $n \in P$ . If n+1 is even then one of n+1, n+2 is even so  $n+1 \in P$ . Otherwise n must be even, so n+2 is also even and again  $n+1 \in P$ .

Example in how to use well-ordering:

THEOREM 7. (Division with remainder) Let  $n, a \in \mathbb{Z}$  with a > 0. Then there are unique  $q, r \in \mathbb{Z}$  with  $0 \le r < a$  such that

$$n = qa + r$$
.

PROOF. Let  $T = \{m \in \mathbb{N} \mid \exists k \in \mathbb{Z} : m = n - ka\}$ . In other words, T is the set of all natural numbers which differ from n by a multiple of a. T is non-empty, since by taking k sufficiently negative we can make n - ka as large as we want (e.g. take k = -|n|). By the well-ordering principle there is  $r = \min T$ . By definition of T, we have  $0 \le r$  and there is  $q \in \mathbb{Z}$  such that r = n - qa. Assume that  $r \ge a$ . Then  $r - a \ge 0$  and r - a = n - (q + 1)a so  $r - a \in T$ , a contradiction. It follows that n = qa + r for some q, r as claimed.

Assume next that also n = q'a + r'. Then

$$q'a + r' = n = qa + r.$$

Assume first that r' > r. Then

$$r'-r=(q-q')a,$$

so by the Lemma,  $a > r' \ge r' - r \ge a$ , a contradiction. By symmetry we can't have r > r' either, so r = r'. It follows that (q - q')a = 0 and since  $a \ne 0$  this means q = q'.

COROLLARY 8. An integer n is odd iff it can be written in the form n = 2k + 1 for some  $k \in \mathbb{Z}$ .

# Math 312: Problem set 1 (due 13/5/11)

# The factorial function and binomial coefficients

Recall that the factorial function is defined by 0! = 1 and for  $n \ge 0$  by  $(n+1)! = (n+1) \cdot n!$ . The *binomial coefficients* are defined for  $0 \le k \le n$  by

$$\binom{n}{k} = \frac{n!}{k!(n-k)!}.$$

If  $k > n \ge 0$  we set  $\binom{n}{k} = 0$  (for example,  $\binom{4}{2} = 6$  while  $\binom{2}{4} = 0$ ).

- 1. For  $0 \le k < n$  show that  $\binom{n}{k+1} + \binom{n}{k} = \binom{n+1}{k+1}$  by a direct calculation. OPT: show that this holds even if  $0 \le n \le k$ .
- 2. For  $n \ge 0$  show that  $\sum_{j=0}^{n} {j \choose 0} = {n+1 \choose 1}$ . *Hint:* once you unwind the definitions of both sides this is not hard.

#### Induction

Use mathematical induction to prove the following assertions:

- 3. Among every three consecutive positive integers there is one that is divisible by 3.
- 4. Fix  $k \ge 0$  and show by induction on n that for  $n \ge 1$ ,  $\sum_{j=0}^{n-1} {j \choose k} = {n \choose k+1}$ . *Hint:* Use the conclusion of problem 1, including the optional part
- 5. (Summation formulas)
  - (a) Show that  $j^2 = 2\binom{j}{2} + \binom{j}{1}$ . This means that  $\sum_{j=0}^{n} j^2 = 2\sum_{j=0}^{n} \binom{j}{2} + \sum_{j=0}^{n} \binom{j}{1}$  (why?). Use problem 4 to establish a formula for  $\sum_{j=0}^{n} j^2$ . *Hint:* You can check your formula (but not the proof) using §1.3.E7.

(b) Express  $j^3$  as a combination of  $\binom{j}{3}$ ,  $\binom{j}{2}$ ,  $\binom{j}{1}$  and use problem 4 to prove a formula for

Hint: Check your formula against §1.3.E8.

# **Divisibility**

An integer a is said to divide the integer b if there is a third integer c such that ac = b. For example, 2 divides 6 since  $2 \cdot 3 = 6$ , but 5 does not divide 6.

- 6. For each integer  $n \in \{6, 12, 17\}$ :
  - (a) List the positive integers which divide n.
  - (b) Find the sum of the divisors of n which are different from n (that is, for each n add all the numbers you got in part (a) except for *n* itself).
  - (c) Is *n* abundant (the sum is bigger than *n*), deficient (the sum is less than *n*) or perfect (the sum is equal to n)?
- 7. Using the lists of divisors from the previous problem:
  - (a) What is the largest number that divides both 6 and 12?
  - (b) What is the largest number that divides both 12 and 17?

REMARK. Perfect numbers are rare and only finitely many are known. It is believed that there are infinitely many even perfect numbers, but this is not known. It is not known if there exist any odd perfect numbers.

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# **Factorials and primes**

- 8. For  $2 \le j \le n$ . Show that n! + j is not prime. Conclude that there are arbitrarily large gaps between consecutive primes.
- 9. For which prime numbers p is p+1 also prime?

REMARK. It is believed (the "twin prime conjecture") that there are infinitely many primes p for which p+2 is also prime.

1.2. Divisibility and the GCD

# 1.2.1. Divisibility.

DEFINITION 9. Let  $a, b \in \mathbb{Z}$ . We say a divides b (and that b is a multiple of a) if there is  $c \in \mathbb{Z}$  such that b = ac. When this holds we write a|b. Otherwise we say a does not divide b and write  $a \nmid b$ .

REMARK 10. Another way to phrase a|b is that the equation ax = b has a solution in  $\mathbb{Z}$ .

NOTATION 11. If a|b with  $a \neq 0$  we write  $\frac{b}{a}$  for the (unique) integer x such that  $a \cdot x = b$ . Note that if  $a \nmid b$  we don't give  $\frac{b}{a}$  any meaning.

EXAMPLE 12. 1|b for all b. 0|b iff b = 0. 15|120. For any  $a, b \in \mathbb{Z}$  we have  $(a - b)|(a^2 - b^2)$ . In particular,  $2^{2^n} - 1 |2^{2^{n+1}} - 1$ .

LEMMA 13. Let  $a \neq 0$  and let a|b. Then  $|b| \geq |a|$ .

PROOF. If 
$$b = ac$$
 we have  $|b| = |a| |c|$ . Also,  $|c| \ge 1$  since  $c \ne 0$  so  $|a| |c| \ge |a|$ .

LEMMA 14. (Euclid) If a divides b and c then a divides  $b \pm c$ .

PROOF. We have  $a\left(\frac{b}{a}\pm\frac{c}{a}\right)=\left(a\cdot\frac{b}{a}\right)\pm\left(a\cdot\frac{c}{a}\right)=b\pm c$  so the equations  $a\cdot x=b\pm c$  have an integer solution.

COROLLARY 15. Let  $n \ge 1$ . Then the only positive common divisor of n, n+1 is 1.

LEMMA 16. If a|b and b|c then a|c.

PROOF. By the associative law, 
$$a \cdot \left(\frac{b}{a} \cdot \frac{c}{b}\right) = \left(a \cdot \frac{b}{a}\right) \cdot \frac{c}{b} = b \cdot \frac{c}{b} = b$$
.

LEMMA 17. (Units) a|b iff (-a)|b.

Because of this, we will only consider *positive* divisors.

### First problem of factorization.

PROBLEM 18. Find all divisors of a given integer.

This turns out to be really hard. We don't know an efficient way to do this.

**1.2.2.** The GCD, Two integers. Let  $a, b \in \mathbb{Z}$  be non-zero. Let D be the set of common divisors of a and b (non-empty since  $1 \in D$ ). D is bounded since every divisor of a is no larger than |a|. Let M be the set of positive common multiples of a, b (non-empty since  $|ab| = \pm ab \in M$ ).

DEFINITION 19.  $(a,b) \stackrel{\text{def}}{=} \gcd\{a,b\} = \max D$ ;  $[a,b] \stackrel{\text{def}}{=} \operatorname{lcm}\{a,b\} = \min M$ . Also, for all  $a \in \mathbb{Z}$ , set (a,0) = a and [a,0] = 0.

FACT 20. Every common divisor of a, b divides (a, b). Every common multiple of a, b is divisible by [a, b].

PROBLEM 21. Given a, b find (a, b) and [a, b].

ALGORITHM 22. (Naive) Try all elements of the finite sets D, M.

Entirely impractical since *finding the divisors* is hard. Euclid discussed a much better idea:

LEMMA 23. (Euclid) Let 
$$a, b \in \mathbb{Z}$$
. Then  $(a, b) = (a - b, b)$ .

PROOF. We prove that both pairs have the same set of common divisors. Indeed, let d divide b. If d also divides a then By Lemma 14 d divides a - b. Conversely, if d divides a - b then by that Lemma d divides a = (a - b) + b.

Since (a,0) = a for all a, and since changing the signs of a,b does not change their gcd (why?) we get a method for calculating the gcd of any two integers. For example:

$$(24,-153) = (153,24)$$

$$= (129,24)$$

$$= (105,24)$$

$$= (81,24)$$

$$= (57,24)$$

$$= (33,24)$$

$$= (24,9)$$

$$= (15,9)$$

$$= (9,6)$$

$$= (6,3)$$

$$= (3,3)$$

$$= (3,0)$$

$$= 3.$$

ALGORITHM 24. (Euclid) Given two integers a,b, output their gcd:

- (1) Replace a with |a|, b with |b|.
- (2) If a < b exchange a and b.
- (3) If b = 0, terminate and output a.
- (4) Else, replace a with a b and go to step 2.

THEOREM 25. The algorithm terminates after finitely many steps and outputs the gcd of (a,b).

PROOF. Consider the changes in the quantity |a|+|b| during the course of the algorithm. Every time we reach step 4, we know that  $a \ge b > 0$ . It follows that at the conclusion of step 4, the quantity has decreased by at least  $b \ge 1$ . Since there is no infinite strictly decreasing sequence of natural numbers (well-ordering), we can reach step 4 only finitely many times. In particular, at some point b = 0 and we terminate. Finally, by Lemma 23, the replacements and exchanges never change the gcd of the two numbers.

In fact, more can be said.

CLAIM 26. (Bezout) Every intermediate value considered by Euclid's Algorithm is of the form xa + yb for some  $x, y \in \mathbb{Z}$ .

PROOF. We prove this by induction on the steps of the algorithm. Certainly this is true at the start, and also changing signs and exchanging a,b doesn't matter. Now assume that at the nth time

we reach step 3, we are looking at the numbers  $a' = xa + yb > b' = za + wb \ge 0$ , where a, b are the initial values and  $x, y, z, w \in \mathbb{Z}$ . At step 4 we will then replace a' with

$$a' - b' = (x - z)a + (y - w)b$$

which is indeed also of this form, so the situation will hold when we reach step 3 for the (n+1)st time.

We have thus proven (by algorithm) the following fact:

THEOREM 27. (Bezout) Given  $a, b \in \mathbb{Z}$  the exist  $x, y \in \mathbb{Z}$  such that (a, b) = xa + yb.

Bezout's theorem admits a direct proof:

PROOF. If a = b = 0 there is nothing to prove, so we assume that at least one of a, b is non-zero, and let  $I = \{ax + by \mid a, b \in \mathbb{Z}\}$ . Note that I is closed under addition and under multiplication by elements of  $\mathbb{Z}$ :  $(ax + by) + q(cx + dy) = (a + qc)x + (b + qd)y \in I$ .

By assumption I contains positive numbers (at least one of |a|, |b| is positive), so let m be the smallest positive element of I. Every common divisor of a,b divides every element of I; in particular  $(a,b) \mid m$ . Conversely, we prove that m divides every element of I. Since  $a,b \in I$  it will follow that m is a common divisor of a,b, hence the greatest common divisor. Let  $n \in I$ . Dividing with remainder (Theorem 7), we can write n = qm + r for some  $0 \le r < m$  and  $q \in \mathbb{Z}$ . Then

$$r = n - qm \in I$$
.

It must be the case that r = 0 (else we'd have a positive member of I smaller than m). Then n = qm and m divides n.

COROLLARY 28. Every common divisor of a,b divides their GCD.

PROOF. Let d divide both a,b. Then for any  $x,y \in \mathbb{Z}$  d|xa and d|yb so d|xa+yb. Now choose x,y so that the xa+yb=(a,b).

#### 1.2.3. The LCM.

DEFINITION 29. Say a, b are relatively prime if (a, b) = 1.

PROPOSITION 30. If a, b are relatively prime then [a,b] = |ab|.

PROOF. By Bezout's Theorem there exist x, y such that xa + yb = 1. Say that az is also a multiple of b. Then  $z = z \cdot 1 = z(xa + yb) = x(az) + (zy)b$  so z is a multiple of b. It follows that az is a multiple of ab, so |ab| is the least positive common multiple.

LEMMA 31. [da, db] = d[a, b].

PROOF. If m is a common multiple of a, b then dm is a common multiple of da, db. Conversely, if m is a common multiple of da, db then m is divisible by d, and  $\frac{m}{d}$  is a common multiple of a, b.

THEOREM 32. Let a, b be non-zero. Then (a, b)[a, b].

PROOF. We have 
$$[a,b] = \left[ (a,b) \frac{a}{(a,b)}, (a,b) \frac{b}{(a,b)} \right] = (a,b) \left[ \frac{a}{(a,b)}, \frac{b}{(a,b)} \right] = (a,b) \frac{ab}{(a,b)^2}.$$

#### 1.2.4. Sets of integers.

DEFINITION 33. Let S be non-empty finite set of integers. We say that  $a \in \mathbb{Z}$  is a *common divisor* of S if a divides every member of S. We say that  $a \in \mathbb{Z}$  is a *common multiple* of S if it is a multiple of every element of S.

EXAMPLE 34. For any non-empty S, 1 is a common divisor for the elements of S and the products of the elements of S is a multiple of all of them.

DEFINITION 35. Assume that S is finite and has a non-zero member. The *greatest common divisor* of S, written gcd(S) is the largest integer which is a common divisor of S. The *least common multiple* of S, written lcm(S) is the smallest positive integer which is a multiple of all elements of S.

NOTATION 36. If  $a_1, \ldots, a_k$  are integers we also write  $(a_1, \ldots, a_k)$  for their GCD and  $[a_1, \ldots, a_k]$  for their LCM.

Note that every common divisor of  $\{a_1, \ldots, a_k\}$  is at most  $|a_1|$ , so only finitely many integers can be the GCD. Similarly, the least common multiple is somewhere between zero and  $\prod_{i=1}^{k} |a_i|$ .

LEMMA 37. 
$$(a_1, \ldots, a_{k+1}) = ((a_1, \ldots, a_k), a_{k+1}).$$

PROOF. Let d be a common divisor of  $a_{k+1}$  and  $(a_1,\ldots,a_k)$ . Then d divides each of  $a_1,\ldots,a_k$  (it divides a common divisor of theirs) so  $d\mid (a_1,\ldots,a_{k+1})$ . It follows that  $((a_1,\ldots,a_k),a_{k+1})\mid (a_1,\ldots,a_{k+1})$ . Conversely, let  $d=(a_1,\ldots,a_{k+1})$ . Then d is a common divisor of  $a_1,\ldots,a_k$  so  $d\mid (a_1,\ldots,a_k)$ . Also,  $d\mid a_{k+1}$ . It follows that d is a common divisor of both  $(a_1,\ldots,a_k),a_{k+1}$  and hence that d divides their GCD, that is that  $(a_1,\ldots,a_{k+1})\mid ((a_1,\ldots,a_k),a_{k+1})$ . Now two positive integers that divide each other are equal.

COROLLARY 38. Algorithm to find the GCD of a list of numbers.

#### 1.2.5. Linear Diophantine equations.

THEOREM 39. The set of integral solutions to ax + by = c is as follows:

- (1) If a = b = 0 then the set is empty if  $c \neq 0$ , all of  $\mathbb{Z}^2$  if c = 0.
- (2) Otherwise, let  $d = \gcd(a, b)$ . Then
  - (a) If  $d \nmid c$  the set is empty.
  - (b) If  $d \mid c$ , let s,t be such that as+bt=d. Then the set of solutions is  $\left\{\left(\frac{sc}{d}+\frac{b}{d}z,\frac{tc}{d}-\frac{a}{d}z\right)\right\}_{z\in\mathbb{Z}}$ .

REMARK 40. Note that "solving an equation" requires doing two tings: showing that every solution is in the set, and showing that every member of the set is a solution.

#### 1.3. Primes

#### 1.3.1. Irreducibles.

DEFINITION 41. Call  $p \in \mathbb{Z}_{>1}$  prime if in every factorization p = ab, one of a, b is 1.

EXAMPLE 42. 1, 2, 3, 5 are irreducible.  $4 = 2 \times 2$  isn't.

THEOREM 43. Every positive integer can be written as a product of primes.

PROOF. Let n be the smallest positive integer which cannot be written as a product of positive primes. Then n itself is not irreducible (nor 1), so n = ab with 1 < a, b < n. But then both a, b are product of irreducibles, hence so is n.

EXAMPLE 44.  $60 = 5 \cdot 12 = 5 \cdot 4 \cdot 3 = 5 \cdot 2 \cdot 2 \cdot 3$ .

THEOREM 45. (Euclid) There are infinitely many primes.

PROOF. Consider P+1 where P is the product of all primes.

REMARK 46. This only shows that there are about  $C \log \log x$  primes up to x.

LEMMA 47. If  $n \ge 1$  is reducible, it has a proper factor  $\le \sqrt{n}$ .

PROOF. This is true about any non-trivial factorization.

COROLLARY 48. Factorization (and primality testing) by trial division.

Example 49.  $126 = 2 \cdot 63 = 2 \cdot 3 \cdot 21 = 2 \cdot 3 \cdot 3 \cdot 7$ .

#### 1.3.2. Primes.

THEOREM 50. Let  $p \in \mathbb{Z}_{>1}$  be prime. Then if p|ab then p divides at least one of a,b.

REMARK 51. This is equivalent to the implication: for all a, b we have  $p \nmid a, p \nmid b \Rightarrow p \nmid ab$ . It is more natural to remove the requirement that p be positive, but it is not traditional to do so.

EXAMPLE 52. 2 is prime.

PROOF. Assume that a, b are odd. Then there are  $k, l \in \mathbb{Z}$  such that a = 2k + 1, b = 2l + 1 so ab = (2k+1)(2l+1) = 2(2kl+k+l) + 1 is also odd.

EXAMPLE 53. 3 is prime.

PROOF. Let  $a, b \in \mathbb{Z}$  not be divisible by 3. Then a = 3q + r and b = 3q' + r' for some  $q, q' \in \mathbb{Z}$  and  $r, r' \in \mathbb{Z}$  with  $1 \le r, r' < 3$ . Then  $r, r' \in \{1, 2\}$ . Since ab = 3(3qq' + r'q + rq') + rr' we have  $3 \mid ab$  iff  $3 \mid rr'$ . But rr' equals one of 1,2,4 and neither is divisible by 3.

PROOF OF THEOREM 50. Conversely, let p be a prime which divides the product ab. Assume that  $p \nmid a$  and consider the GCD d = (a, p). Since  $d \mid p$ , d = 1 or d = p. Since p does not divide a,  $d \neq p$  so d = 1. It follows that there exist x, y such that xp + ya = 1, at which point

$$b = xpb + yab$$
.

Then p divides b since it divides both xpb and yab.

REMARK 54. Conversely, let p be a number such that p|ab implies p|a or p|b. Then either p is a unit or prime. Indeed, if p = ab then p divides one of the factors, so without loss of generality we have  $p \mid a$ . This forces  $|p| \le |a|$  and hence  $|b| \le 1$ . It follows that |b| = 1 so b is a unit.

THEOREM 55. ("Fundamental Theorem of Arithmetic") Every positive integer has a factorization into primes, unique up to reordering the factors.

PROOF. What we need to show is: assume that  $n = \prod_{i=1}^{I} p_i = \prod_{j=1}^{J} q_j$  with  $\{p_i\}, \{q_j\}$  primes. Then I = J and there is a permutation  $\pi \in S_n$  for which  $p_i = q_{\pi(i)}$ . Let n be the smallest positive integer for which this fails. Let p be a prime divisor of n. Then there is i such that  $p_i = p$  and j such that  $q_j = p$  ...

NOTATION 56. We may uniquely write every non-zero integer in the form  $n = \varepsilon \prod_p p^{e_p}$  where  $\varepsilon \in \{\pm 1\}$  and  $e_p$  are non-negative integers, equal to zero for all but finitely many p.

Example 57. 
$$v_p(n!) = \left\lfloor \frac{n}{p} \right\rfloor + \left\lfloor \frac{n}{p^2} \right\rfloor + \left\lfloor \frac{n}{p^3} \right\rfloor + \cdots$$

PROPOSITION 58. Every positive divisor of  $\prod_p p^{e_p}$  is of the form  $\prod_p p^{f_p}$  where  $0 \le f_p \le e_p$ .

Theorem 59.  $(\prod_{p} p^{e_p}, \prod_{p} p^{f_p}) = \prod_{p} p^{\min\{e_p, f_p\}}$  while  $[\prod_{p} p^{e_p}, \prod_{p} p^{f_p}] = \prod_{p} p^{\max\{e_p, f_p\}}$ .

COROLLARY 60. (a,b)[a,b] = ab.

PROOF. 
$$\min\{e_p, f_p\} + \max\{e_p, f_p\} = e_p + f_p$$
.

# 1.3.3. Distribution of primes (not examinable).

DEFINITION 61.  $\pi(x) = \#\{1 \le p \le x \mid p \text{ prime}\}.$ 

CONJECTURE 62. (Gauss circa 1800)  $\lim_{x\to\infty}\pi(x)\cdot\frac{\log}{x}=1$ . More precisely,  $\pi(x)\sim \operatorname{Li}(x)\stackrel{def}{=}\int_2^x\frac{dt}{\log t}$  ("integers are prime with probability  $\frac{1}{\log t}$ ")

THEOREM 63. (Chebychev 1850) For all large enough x,

$$0.9 \frac{x}{\log x} \le \pi(x) \le 1.1 \frac{x}{\log x}.$$

THEOREM 64. (de la Vallée-Poussin, Hadamard 1896, following Riemann 1859)  $|\pi(x) - \text{Li}(x)| \le Cxe^{-c\sqrt{\log x}}$ .

CONJECTURE 65. (Riemann 1859)  $|\pi(x) - \text{Li}(x)| < C\sqrt{x}\log x$ .

THEOREM 66. (Dirichlet 1837) Barring local obstruction, every AP contains infinitely many primes.

(Chebotarev) Let  $\pi(q, a; x) = \#\{1 \le p \le x \mid p \equiv a(q) \text{ prime}\}$ . Then  $\left|\pi(x) - \frac{1}{\phi(q)}\operatorname{Li}(x)\right| \le Cxe^{-c\sqrt{\log x}}$ .

Conjecture 67. (ERH) For (a,q)=1,  $\left|\pi(q,a;x)-\frac{1}{\phi(q)}\operatorname{Li}(x)\right|\leq C\sqrt{x}\log x$ .

#### 1.3.4. Special primes.

- Twin primes.
- Fermat numbers:  $F_n = 2^{2^n} + 1$ . Prime for n = 0, 1, 2, 3, 4. Fermat Conj all prime; Euler showed  $641|2^{2^{3^2}} + 1$ . No other Fermat primes known.
  - $(F_n, F_m) = 1$  if  $n \neq m$ , so infinitely many primes.
- Mersenne numbers:  $2^p 1$ . Many known; probably there are  $\infty$  many but it is open.
  - Cole, October 1903 meeting of the AMS:  $2^{67} 1 = 193,707,721 \times 761,838,257,287$ .
  - Euclid:  $2^p 1$  prime then  $2^{p-1}(2^p 1)$  perfect.
  - Euler: converse.

PROPOSITION 68.  $(F_n, F_m) = 1$  if  $n > m \ge 0$ .

PROOF. For any integer x we have

$$(x^2+1,x+1) = ((x+1)^2-2(x+1)+2,x+1) = (2,x+1).$$

In particular, for  $x=2^{2^m}$  which is even we find  $(F_{m+1},F_m)=1$ . Assume now that n>m+1, and for  $0 \le j < 2^{n-m-1}$  write  $F_{n,j}=2^{2^n-j2^{m+1}}+1$  so that  $F_{n,0}=F_n$  and  $F_{n,2^{n-m-1}-1}=2^{2^{m+1}}+1=F_{m+1}$ . Then, for  $0 \le j < 2^{n-m-1}-1$  we have

that is  $(F_{n,j},F_m)=(F_{n,j+1},F_m)$ . It follows by induction that  $(F_n,F_m)=(F_{m+1},F_m)=1$ .

COROLLARY 69. There are infinitely many primes.

PROOF. No prime divides two of the  $F_n$ .

REMARK 70. Note that this proof only produces n primes up to  $2^{2^n} + 1$ , i.e. about  $\log \log x$  primes up to x.

# Math 312: Problem set 2 (due 19/5/11)

# **Prime factorization**

- 1. (§3.5.E2) Find the prime factorization of 111, 111.
- 2. Let (a,c) = 1. Show that (a,bc) = (a,b). *Hint:* Factor a, b, c into primes and calculate both sides explicitly.
- 3. Recall that  $\binom{n}{k} = \frac{n!}{k!(n-k)!}$ .
  - (a) Give (with proof) a finite list of primes which contains all prime divisors of  $\binom{30}{10}$ .
  - (b) For each prime on your list, find the number of times it divides  $\binom{30}{10}$ . Give the prime factorization of this number.
- 4. We consider the equation  $2^x + 3^y = z^2$  for unknown  $x, y, z \in \mathbb{Z}_{>0}$ .
  - (a) (The case y = 0) Find all non-negative integral solutions to  $2^x + 1 = z^2$ . *Hint*: Start by showing that both z-1 and z+1 must be powers of 2.
  - (b) (The case x = 0) Find all non-negative integral solutions to  $1 + 3^y = z^2$ . *Hint*: Which powers of 3 differ by 2?
  - (c) Let (x, y, z) be a solution with both of x, y positive and even. Show that z 2<sup>x/2</sup> = 1. *Hint*: If 3 divides both z 2<sup>x/2</sup> and z + 2<sup>x/2</sup> it would divide their difference.
    (d) Continuing (c), show that 3<sup>y</sup> = 2<sup>1+x/2</sup> + 1 and find all solutions to this equation.
  - *Hint*: Both  $3^{y/2} \pm 1$  must be powers of 2.
  - RMK We will show in future problem sets that if (x, y, z) is a solution to the equation above and x, y are positive then x, y are even.

# **Euclid's Algorithm**

5. For each pair of integers a, b use Bezout's extension of Euclid's algorithm to find gcd(a, b) and integers x, y such that  $ax + by = \gcd(a, b)$ . Give your intermediate calculations.

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- (a) a = 5, b = 2.
- (b) a = 60, b = 36.
- 6. Let a > b > 0.
  - (a) Show that  $(2^a 1, 2^b 1) = (2^a 1, 2^{a-b} 1)$ . Hint: Euclid's Lemma + problem 2.
  - (b) Show that  $(2^a 1, 2^b 1) = 2^{(a,b)} 1$ . *Hint*: Strong induction on a + b.
  - (c) Show that  $(x^a 1, x^b 1) = x^{(a,b)} 1$  for all  $a \ge b \ge 0$  and all  $x \ge 2$ .

### **Primes**

7. For a positive integer n, show that n! + 1 has a prime divisor> n. Conclude that there are infinitely many primes.

For the next two problems use the identities  $x^n - y^n = (x - y) \sum_{k=0}^{n-1} x^k y^{n-1-k}$  and (for n odd)  $x^n + y^n = (x + y) \sum_{k=0}^{n-1} (-1)^k x^k y^{n-1-k}$ .

- 8. Let a, n be integers with  $a \ge 1$ ,  $n \ge 2$  such that  $a^n 1$  is prime.
  - (a) Show that a = 2.
  - (b) Show that n is prime.

*Hint:* This follows from 6(b) or from the identities above.

9. Let a, b, n be positive integers (with ab > 1) such that  $a^n + b^n$  is prime. Show that n is a power of 2.

*Hint*: Try ruling out n = 6 before tackling the general case.

# **Linear equations**

- 10. We study the equation ax + by = 0 where not both a, b are zero.
  - (a) Let a, b be relatively prime. Show that the solutions to ax + by = 0 are precisely the pairs of the form x = bz, y = -az with  $z \in \mathbb{Z}$  arbitrary.

*Hint*: Note that you both need to verify that these are solutions and to show that every solution is of this form.

- (b) Now let d = (a,b) be anything. Find all solutions to the equation. *Hint:* Divide by d.
- (c) Use part (a) to find all solutions to 5x + 6y = 1. *Hint*: 6 - 5 = 1.

# Supplementary problems (not for submission): A counting proof of the infinitude of primes

- A. In the factorization  $n = \prod_p p^{e_p}$  show that  $e_p \le \log_2 n$ .
- B. Assume that  $\{p_j\}_{j=1}^r$  is the set of all primes, and let  $x \ge 2$ . Show that there are at most  $(1 + \log_2 x)^r$  integers between 1 and x.
- C. Show that  $\frac{(1+\log_2 x)^r}{x} \to 0$  as  $x \to \infty$  and derive a contradiction.
- D. Use this idea to show that  $\pi(x) \ge C \frac{\log_2 x}{\log_2 \log_2 x}$ .
- E. (unrelated) Let  $n = \prod_p p^{e_p} \ge 1$ . Show that n has  $\tau(n) = \prod_p (e_p + 1)$  positive divisors.

#### CHAPTER 2

# Congruences

### 2.1. Arithmetic in congruences

DEFINITION 71. For  $a, b, m \in \mathbb{Z}$  with  $m \ge 2$ ,  $a \equiv b(m)$  iff  $m \mid a - b$ .

PROPOSITION 72. Congruence mod m is an equivalence relation.

LEMMA 73.  $a \equiv b(m)$  iff b = a + km for some  $k \in \mathbb{Z}$ .

THEOREM 74. (Arithmetic) If  $a \equiv a', b \equiv b'(m)$  then

$$a \pm b \equiv a' + b'(m)$$

and

$$ab \equiv a'b'(m)$$
.

PROOF. Say a'=a+km, b'=b+lm. Then  $a'\pm b'=a\pm b+(k\pm l)m$  and a'b'=ab+(al+bk+kl)m.

FACT 75. (Division Thm) Every congruence class mod m contains a unique representative  $0 \le r < m$ .

DEFINITION 76. The *reduction* of a mod m is that  $0 \le r < m$  which is congruent to a.

EXAMPLE 77. (Divisibility tests)  $m \mid a$  iff the reduction of  $a \mod m$  is zero.

LEMMA 78. For all  $k \ge 0$ ,  $10^k \equiv 1 (9)$  (proof by induction)

COROLLARY 79. (Divisibility test) For all  $n \ge 0$ ,  $n \equiv S(n)(9)$  where S(n) is the sum of digits of n.

DEFINITION 80. Say that b is a modular inverse of  $a \pmod{m}$  if  $ab \equiv 1 \pmod{m}$ . Say that a is invertible if it has a modular inverse.

THEOREM 81. *a is invertible*  $\mod m$  *iff* (a,m) = 1.

PROOF. Assume that  $ab \equiv 1$  (m). Then 1 = ab + km. Thus any prime dividing both a and m must divide 1. Conversely, let a, m be relatively prime. By Bezout's Theorem (Theorem 27), there are x, y such that ax + my = 1, which means that  $ax \equiv 1$  (m).

Example 82.  $2^7 \equiv -\bar{5}$  (641) while  $2^4 \equiv -5^4$  so  $2^{32} + 1 = 2^4 \left(2^7\right)^4 + 1 \equiv -5^4 \left(-\bar{5}\right)^4 + 1 \equiv -1 \left(5 \cdot \bar{5}\right)^4 + 1 \equiv 0$  (641).

THEOREM 83. (CRT) Let  $\{m_j\}_{j=1}^J$  be pairwise relatively prime, and let  $M = \prod_j m_j$ . Let  $\{a_j\}_{j=1}^J \subset \mathbb{Z}$ . Then there is  $a \in \mathbb{Z}$ , unique mod M, such that  $a \equiv a_j(m_j)$ .

PROOF. Assume first that  $a_1 = 1$  and that  $a_j = 0$  for  $2 \le j \le J$ . Let  $N = \prod_{j=2}^J m_j$  so that  $M = m_1 N$ . It is then enough to find a such that  $a \equiv 1 \, (m_1)$  while  $a \equiv 0 \, (N)$ . Since  $(N, m_1) = 1$  can take  $a = N\bar{N}$  where  $\bar{N}$  is any inverse of  $N \mod m_1$ . It follows that there exist  $\{y_j\}_{j=1}^J$  such that  $y_i \equiv \delta_{ij} \, (m_j)$ . For existence set  $a = \sum_j y_j a_j$ . For uniqueness by subtraction it is enough to consider the case  $a \equiv 0 \, (m_j)$  for all j, for which a must be divisible by M.

EXAMPLE 84. Divide the six residue classes mod 6 into their classes mod 2 and 3.

EXAMPLE 85. Solve  $x \equiv 1(3)$ ,  $x \equiv 2(5)$  and  $x \equiv 3(7)$ . Indeed  $35 \cdot 2 = 70 \equiv 1(3)$ ,  $21 \equiv 1(5)$  and  $15 \equiv 1(7)$ . It follows that the solution is  $x \equiv 70 \cdot 1 + 21 \cdot 2 + 15 \cdot 3(3 \cdot 5 \cdot 7)$ , that is  $x \equiv 157 \equiv 52(105)$ .

# Math 312: Problem set 3 (due 25/5/11)

# Calculation

- 1. (Dec 2010 final exam) Let  $n = 3^{100}$  and let a be the 2n-digit number 858585...8585 where the digits 85 repeat n times. Is a divisible by 9? Prove your answer.
- 2. (Dec 2005 final exam)
  - (a) Show that  $3^6 \equiv 1(7)$

*Hint:* Calculate  $3^2$  or  $3^3$  mod 7 first.

(b) Let  $a \equiv b$  (6). Show that  $3^a \equiv 3^b$  (7).

*Hint*: What can you say about  $3^{|a-b|}$ ? Problem 8 may be useful.

- (c) Today is Thursday. What day will it be  $10^{200,000,000,000}$  days from now?
- 3. ( $\S4.1.E20$ ) Find the least non-negative residue mod 13 of the following numbers: 22, -100, 1001.
- 4. (squares mod small numbers)
  - (a) For each m = 3,4 find all residues  $0 \le a < m$  which are square mod m (in other words for which there is an integer solution to  $x^2 \equiv a(m)$ ).

*Hint*: Let x range over the residues mod m and see what values of a you get.

- (b) Find an integer x such that  $x^2 \equiv -1(5)$ .
- 5. Find all solutions to:  $15x \equiv 9(25)$ ; also to  $2x + 4y \equiv 6(8)$ .
- 6. (CRT)
  - (a) (§4.3.E10) Find an integer that leaves a remainder of 9 when divided by 10 or 11 but is divisible by 13.
  - (b) (§4.3.E12) If eggs are removed from a basket 2,3,4,5,6 at a time, 1,2,3,4,5 eggs remain, respectively. If eggs are removed 7 at a time, no eggs remain. What is the least possible number of eggs in the basket?

*Hint*: Note that -1 satisfies the congruence conditions modulu 2,3,4,5,6 hence mod their LCM.

#### **Problems**

- 7. Powers and irrationals
  - (a) Let  $n = \prod_p p^{e_p}$  be the prime factorization of a positive integer and let  $k \ge 2$ . Show that in the prime factorization of  $n^k$  every exponent is divisible by k. Conversely, let  $m = \prod_p p^{f_p}$  where  $k|f_p$  for all p. Show that m is the kth power of a positive integer.
  - (b) Show that  $\sqrt{2}$  is not an integer, that is that there is no integer solution to  $x^2 = 2$ . Hint: What is the exponent of 2 in the prime factorization of 2? What do you know about the exponent of 2 in the prime factorization of  $x^2$ ?
  - (c) Show that  $\sqrt{2}$  is not a rational number, that is that there are no positive integers x, y such that  $\left(\frac{x}{y}\right)^2 = 2$ .

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*Hint*: Consider the exponent of 2 on both sides of  $x^2 = 2y^2$ .

SUPP Show that  $\sqrt{2} + \sqrt{3}$  is irrational.

*Hint*: Squaring shows that if this number is rational then so is  $\sqrt{6}$ ...

- 8. Let  $a \equiv b(m)$ . Show that  $a^n \equiv b^n(m)$  for all n > 0.
- 9. Consider the numbers  $2^x \mod 3$  and  $3^y \mod 4$ .
  - (a) Let  $2^x + 3^y = z^2$  for some integers  $x, y, z \ge 0$  where  $x, y \ge 1$ . Show that  $(-1)^x \equiv z^2(3)$ .
  - (b) Use problem 4 to show that  $(-1)^x \equiv z^2(3)$  forces x to be even. Hint: Is (-1) a square mod 3?
  - (c) Now show that  $(-1)^y \equiv z^2(4)$ .
  - (d) Finally, show that this forces y to be even.
- 10. For  $n = \sum_{j=0}^{J} 10^j a_j$  set  $T(n) = \sum_{j=0}^{J} (-1)^j a_j$  (i.e. add the even digits and subtract the odd digits).
  - (a) Show that  $T(n) \equiv n(11)$ .
  - (b) Is the number from problem 1 divisible by 11? Justify your answer.
- 11. (Gaps between squarefree numbers)
  - (a) Let  $\{p_j\}_{j=1}^J$  be distinct primes. Show that there exist positive integers x such that for all  $1 \le j \le J$ ,  $p_j^2 | x + j$ .

*Hint*: Rewrite the condition as a congruence condition on x and apply the CRT.

(\*b) Call a number "squarefree" if it is not divisible by the square of a prime (15 is squarefree but 45 isn't). Show that there are arbitrarily large gaps between squarefree numbers.

# Supplementary problems (not for submission)

- A. Show that every non-zero rational number can be uniquely written in the form  $\varepsilon \prod_p p^{e_p}$  where  $\varepsilon \in \{\pm 1\}$ ,  $e_p \in \mathbb{Z}$  and  $\{p \mid e_p \neq 0\}$  is finite. Show that a rational number is a kth power iff  $\varepsilon$  is a kth power and  $k|e_p$  for all p.
- B. (The *p*-adic norm) For a rational number  $a = \varepsilon \prod_p p^{e_p}$  with a factorization as above set  $|n|_p = p^{-e_p}$  (and  $|0|_p = 0$ ).
  - (a) Show that  $|a+b|_p \le \max\{|a|_p,|b|_p\} \le |a|_p + |b|_p$  and  $|ab|_p = |a|_p |b|_p$ .
  - (b) Show that  $d_p(a,b) + d_p(b,c) \ge d_p(a,c)$ .

### 2.2. Linear Congruences

THEOREM 86. Consider the equation  $ax \equiv b(m)$ , and let d = (a,m). If  $d \nmid b$  there are no solutions. Otherwise it is equivalent to  $\frac{a}{d}x \equiv \frac{b}{d}(\frac{m}{d})$ . If (a,m) = 1 then  $ax \equiv b$  iff  $x \equiv \bar{a}b$ .

# 2.3. The multiplicative group

2.3.1. Multiplicative order, Euler's Theorem and Fermat's Little Theorem. Examine powers of 2 mod 7, mod 11. Powers of 5 mod 11 and see periodicity. Examine powers of 2 mod 6 and see periodicity but no 1.

DEFINITION 87. For (a, m) = 1 the multiplicative order of a mod m is  $\operatorname{ord}_m(a) = \min \{n \ge 1 \mid a^n \equiv 1 \ (m) \}$ .

PROPOSITION 88. Let (a, m) = 1. Then  $a^r \equiv a^s(m)$  iff  $r \equiv s(\operatorname{ord}_m(a))$ .

#### **2.3.2. Wilson's Theorem.** We evaluate the

#### 2.3.3. Examples.

- Pollard's p-1 method.
- The order of 2 mod  $2^n 1$  is n: it divides n since  $2^n \equiv 1(2^n 1)$  but can't be smaller. Factor  $3^{20} + 3^{19} 12 = 12(3^{18} 1) = 12(3^9 + 1)(3^9 1)$ .
- See that 35 is not prime:  $2^{34} \not\equiv 1(35)$ .

ALGORITHM 89. Modular exponentiation by repeated squaring

### 2.4. Primality testing

Pseudoprimes, Charmichael numbers, strong pseudoprimes,

# Math 312: Problem Set 4 (due 1/6/11)

### Linear equations

1. §5.5.E12.

#### **Multiplicative Order**

- 2. Let *n* be a pseudoprime to base 2. Show that  $m = 2^n 1$  is also a pseudoprime to base 2. *Hint*: Show that n|m-1 and use the fact that you know the order of 2 mod *m*.
- \*3. Let *p* be a prime divisor of the *n*th Fermat number  $F_n = 2^{2^n} + 1$ .
  - (a) Find the order of  $2 \mod p$ .
  - (b) Show that  $p \equiv 1 (2^{n+1})$ .
  - (c) Show that for any  $a \ge 1$  there are infinitely many primes p for which the order of  $2 \mod p$  is divisible by  $2^a$ .
  - RMK Note that (b) simplifies the search for prime divisors of Fermat numbers. We will later show that  $p \equiv 1 (2^{n+2})$  holds.
- 4. Elements of order  $2 \mod m$ .
  - (a) Let p be odd, and let  $k \ge 1$ . Show that the congruence  $x^2 \equiv 1$  ( $p^k$ ) has only the two obvious solutions  $x \equiv \pm 1$  ( $p^k$ ).

*Hint*: Can both x - 1, x + 1 be powers of p?

- (\*b) Let n be an odd number, divisible by exactly r distinct primes. Set up a bijection between congruence classes mod n satisfying  $x^2 \equiv 1$  (n) and functions  $f \in \{\pm 1\}^r$ . Conclude that there are precisely  $2^r$  congruence classes mod n which solve the equation.
- 5. Using Fermat's Little Theorem, show that for all integers n,  $30|n^9 n$ . *Hint*: For each prime p|30 show that  $n^p n|n^9 n$  as polynomials.

#### Wilson's Theorem

- 6. We will show that if n > 6 is composite then  $(n-1)! \equiv 0(n)$ .
  - (a) (The easy case) Assume first that n is divisible by at least two distinct primes, that is that  $n = \prod_{j=1}^r p_j^{k_j}$  for some distinct primes  $p_j$  where  $k_j \ge 1$  for all j and  $r \ge 2$ . Show that  $(n-1)! \equiv 0$  (n).

*Hint*: It is enough to show the congruence mod each  $p_j^{k_j}$  separately. Why is (n-1)! divisible by  $p_j^{k_j}$ ?

- (b) Let p be prime and let  $k \ge 3$ . Show that  $p^k | (p^k 1)!$  *Hint:* Find some powers of p dividing the factorial.
- (c) Let  $p \ge 3$  be prime. Show that  $p^2 | (p^2 1)!$  Hint: Now you need to consider multiples of p as well.

RMK Note that  $3! \neq 0$  (4). Ensure that your solution to (c) used the fact that  $p \neq 2$  at some point!

### The Euler Function and RSA

Recall that  $\phi(m) = \#\{1 \le a \le m \mid (a,m) = 1\}$ , and that for p prime  $\phi(p) = p - 1$ .

- 8. Explicit calculations.
  - (a) Calculate  $\phi(4), \phi(9), \phi(12), \phi(15)$ .
  - (b) Show that  $\phi(12) = \phi(3)\phi(4)$  and  $\phi(15) = \phi(3)\phi(5)$  but that  $\phi(4) \neq \phi(2) \cdot \phi(2)$ ,  $\phi(9) \neq \phi(3) \cdot \phi(3)$ .
- 9. Let p,q be distinct primes and let m = pq.
  - (a) Show that there are p+q-1 integers  $1 \le a \le m$  which are not relatively prime to m. *Hint*: What are the possible values of gcd(a, m)? For which a do they occur?
  - (b) Show that  $\phi(pq) = (p-1)(q-1)$ .

RMK This means in particular that  $\phi(pq) = \phi(p)\phi(q)$ .

(c) Give a formula for p+q in terms of  $m, \phi(m)$ .

SUPP Show how to factor m given m,  $\phi(m)$ .

- 10. Fix an integer m and two positive integers d, e so that  $de \equiv 1$  ( $\phi(m)$ ). Define functions E, D by  $E(x) = x^e \mod m$  and  $D(y) = y^d \mod m$  (in other words, raise to the appropriate power and keep remainder mod m).
  - (a) Let  $M = \{1 \le a \le m \mid (a, m) = 1\}$  be the set of invertible residues  $(\phi(m))$  is the size of this set). Show that both D, E map the set M into itself.
  - (b) Show that for any  $x, y \in M$ , D(E(x)) = x and E(D(y)) = y. *Hint:* Euler's Theorem.

# **Supplementary problems (not for submission)**

A. (The binomial formula) Prove by induction on  $n \ge 0$  that for all x, y,

$$(x+y)^n = \sum_{k=0}^n \binom{n}{k} x^k y^{n-k}.$$

- B. Let p be an odd prime.
  - (a) Show that  $(p-1)! \equiv (-1)^{\frac{p-1}{2}} \left( \left( \frac{p-1}{2} \right)! \right)^2 (p)$ . Conclude that if  $p \equiv 1$  (4) then there is  $a \in \mathbb{Z}$  such that  $a^2 \equiv -1$  (p).
  - (b) Conversely, assume that  $a^2 \equiv -1(p)$  for some integer a. Show that the order of  $a \mod p$  is exactly 4 and conclude that  $p \equiv 1(4)$ .

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C. Let p be a prime and let  $0 \le k < p$ . Show that  $\binom{p-1}{k} \equiv (-1)^k (p)$ .

#### CHAPTER 3

# **Arithmetic functions**

#### 3.1. Dirichlet convolution

DEFINITION 90. An *arithmetical function* is a function  $f: \mathbb{Z}_{\geq 1} \to \mathbb{R}$  (more generally, to  $\mathbb{C}$ ). EXAMPLE 91. Some standard functions.

- All-ones function I(n) = 1, identity function N(n) = n, delta-function  $\delta(n) = \begin{cases} 1 & n = 1 \\ 0 & n \neq 1 \end{cases}$ .
- Characteristic function of the primes:  $P(n) = \begin{cases} 1 & n \text{ prime} \\ 0 & \text{not} \end{cases}$  for which  $\pi(x) = \sum_{n \leq x} P(n)$ .

  The von-Mangoldt function  $\Lambda(n) = \begin{cases} \log p & n = p^k, k \geq 1 \\ 0 & \text{else} \end{cases}$  which is the "right way" to

count primes using *Chebychev's function*  $\psi(x) = \sum_{n \le x} \Lambda(n)$ .

- Let  $n = \prod_{p} p^{e_p}$ . Then  $\omega(n) = \#\{p \text{ prime } : p|n\}, \Omega(n) = \sum_{p} e_p$ . From these get the Möbius function  $\mu(n) = \begin{cases} (-1)^{\omega(n)} & n \text{ squarefree} \\ 0 & \text{else} \end{cases}$ , the Liouville function  $\lambda(n) = (-1)^{\Omega(n)}$ .
- $\phi(n) = \#\{0 \le a < n \mid (a,n) = 1\}, \ \tau(n) = \sum_{d|n} 1, \ \sigma(n) = \sum_{d|n} d, \ \sigma_k(n) = \sum_{d|n} d^k.$

DEFINITION 92. Call f multiplicative if f(mn) = f(m) f(n) if (m,n) = 1.

REMARK 93. This usually has to do with the CRT.

EXAMPLE 94.  $\phi(n)$  is multiplicative (proof later).

LEMMA 95. If f is multiplicative and  $f(1) \neq 1$  then f(n) = 0 for all n.

PROOF. 
$$f(n) = f(n \cdot 1) = f(n)f(1)$$
 so  $(1 - f(1)) \cdot f(n) = 0$  for all  $n$ .

PROPOSITION 96. f multiplicative and  $n = \prod_{p} p^{e_p}$  then  $f(n) = \prod_{p} f(p^{e_p})$ .

PROOF. Induction on number of prime factors of *n*.

Corollary 97. 
$$\phi(p^k) = p^k - p^{k-1} = p^k \left(1 - \frac{1}{p}\right) so \ \phi(\prod_p p^{e_p}) = \prod_{p|n} \left(p^{e_p} - p^{e_p-1}\right) = \prod_{p|n} p^{e_p} \left(1 - \frac{1}{p}\right) = n \prod_{p|n} \left(1 - \frac{1}{p}\right).$$

DEFINITION 98. f is completely multiplicative if f(mn) = f(m)f(n) for all n.

Example 99. I(n),  $\delta(n)$ ,  $N^k(n)$ .

DEFINITION 100. The Dirichlet convolution of f,g is the arithmetical function

$$(f*g)(n) = \sum_{d|n} f(d)g(\frac{n}{d}) = \sum_{de=n} f(d)g(e)$$

The second definition shows that f \* g = g \* f. By convention the sum is over positive divisors and factorizations only.

EXAMPLE 101. Calculate  $(\phi * I)(n)$  for small values of n, note that  $\phi * I = N$ .

THEOREM 102.  $\phi *I = N$ .

PROOF. (Textbook) Combinatorial  $-\sum_{d|n} \phi(\frac{n}{d}) = n$  since this counts the integers between 1, n according to their gcd with n.

THEOREM 103. f, g multiplicative then so if f \* g.

PROOF. Let  $gcd(m_1, m_2) = 1$ . Then there is a bijection between divisors  $d|m_1m_2$  and pairs  $d_1|m_1, d_2|m_2$  by  $(d_1, d_2) \mapsto d_1d_2$  and  $d \mapsto (gcd(d, m_1), gcd(d, m_2))$ . It follows that

$$(f * g)(m_1 m_2) \stackrel{\text{def}}{=} \sum_{d|m_1 m_2} f(d)g(\frac{m_1 m_2}{d})$$

$$= \sum_{d_1|m_1} \sum_{d_2|m_2} f(d_1 d_2)g\left(\frac{m_1}{d_1} \cdot \frac{m_2}{d_2}\right)$$

$$= \sum_{d_1|m_1} \sum_{d_2|m_2} f(d_1)f(d_2)g\left(\frac{m_1}{d_1}\right)g\left(\frac{m_2}{d_2}\right) \qquad f,g \text{ multiplicative}$$

$$= \sum_{d_1|m_1} f(d_1)g\left(\frac{m_1}{d_1}\right) \sum_{d_2|m_2} f(d_2)g\left(\frac{m_2}{d_2}\right)$$

$$= ((f * g)(m_1)) \cdot ((f * g)(m_2)).$$

Example 104.  $\tau = I * I$ ,  $\sigma_k = I * N^k$ .

REMARK 105. f, g completely multiplicative doesn't mean f \* g is so, e.g.  $\tau = I * I$ .

Example 106.  $\tau(p^k) = k+1$  so  $\tau(\prod_p p^{e_p}) = \prod_p (e_p+1)$ .  $\sigma(p^k) = \frac{p^{k+1}-1}{p-1}$  so  $\sigma(\prod_p p^{e_p}) = \prod_{p \mid n} \frac{p^{e_p+1}-1}{p-1}$ .

PROBLEM 107. (Past final)  $\tau(n) = 77$  and 6|n. Find n.

#### 3.2. Mersenne primes and perfect numbers

DEFINITION 108. (cf PS1) Call *n* deficient if  $\sigma(n) < 2n$ , abundant if  $\sigma(n) > 2n$  and perfect if  $\sigma(n) = 2n$ .

EXAMPLE 109. 6,28.

Not clear if odd perfect numbers exist. We'll study even perfect numbers.

- Let  $n = 2^s m$  be perfect with m odd. Then by multiplicativity  $\sigma(n) = \sigma(2^s)\sigma(m) = (2^{s+1}-1)\sigma(m)$ .
  - We used  $(2^s, m) = 1$ . Not enough to say "2<sup>s</sup> even, m odd".
- By assumption also  $\sigma(n) = 2n = 2^{s+1}m$ . It follows that  $2^{s+1} | (2^{s+1} 1)m$ . Since  $(2^{s+1}, 2^{s+1} 1) = 1$  (they are consecutive) we have  $2^{s+1} | \sigma(m)$  so write  $\sigma(m) = 2^{s+1}t$ .

- We then have  $2^{s+1}m = (2^{s+1}-1)2^{s+1}t$  that is  $m = (2^{s+1}-1)t$ .
- If t > 1 then  $1, t, (2^{s+1} 1)t$  are distinct divisors of m so  $\sigma(m) \ge 1 + t + (2^{s+1} 1)t = 1 + 2^{s+1}t > \sigma(m)$ , a contradiction.
- It follows that  $m = 2^{s+1} 1$  and that  $\sigma(m) = 2^{s+1} = m + 1$ . In particular, m has no other divisors that 1, m so m is prime.
- By PS2, if  $2^{s+1} 1$  is prime then s+1 itself is prime. We have shown that every even perfect number is of the form  $2^{p-1}(2^p-1)$  where  $p, 2^p-1$  are both prime.
- Conversely, if  $p, 2^p 1$  are both prime then  $\sigma\left(2^{p-1}(2^p 1)\right) = \sigma\left(2^{p-1}\right)\sigma\left(2^p 1\right) = (2^p 1)(1 + 2^p 1) = 2 \cdot 2^{p-1}(2^p 1)$ .

THEOREM 110. An even number is perfect iff it is of the form  $2^{p-1}(2^p-1)$  for a prime of the form  $2^p-1$ .

DEFINITION 111. The numbers of the form  $M_n = 2^n - 1$  are called *Mersenne numbers*.

PROPOSITION 112. Let q, p be primes with  $q|2^p - 1$ . Then  $q \equiv 1$  (p).

PROOF. We have seen that the order of  $2 \mod 2^p - 1$  is  $p (2^p \equiv 1 (2^p - 1))$  so the order must divide p, but it is not 1). The same argument works mod q: if  $q|2^p - 1$  then  $2^p \equiv 1 (q)$  so  $\operatorname{ord}_q(2)|p$ . But the order is not  $1 (2 \not\equiv 1 (q))$  so it is p. By Fermat's Little Theorem, the order of any number mod q divides q - 1 so p|q - 1.

EXAMPLE 113. The first few Mersenne primes and associated perfect numbers are:

- $2^2 1 = 3$ ;  $2^1 \cdot 3 = 6$
- $2^3 1 = 7$ ;  $2^2 \cdot 7 = 28$
- $2^5 1 = 31$  if not prime would have a prime divisor  $\equiv 1(5)$  and  $< 6 = \sqrt{36}$  which is impossible. The perfect number is  $2^4 \cdot 31 = 496$ .
- $2^7 1 = 127$  if not prime would have a prime divisor  $\equiv 1 \, (7)$  and  $< 12 = \sqrt{144}$  but there are no such primes.  $2^6 \cdot 127 = 8128$ .

EXAMPLE 114.  $2^{11} - 1 = 2047$  – if not prime would have a prime divisor  $\equiv 1 \, (11)$  and  $< 50 = \sqrt{2500}$ . The only prime in this range is 23 and indeed  $\frac{2047}{23} = 89$ .

# **Math 312: Problem Set 5 (due 8/6/11)**

# **Arithmetic functions**

- 1. In this exercise, f, g are multiplicative functions. You are also given that for  $n \in \{2, 3, 4, 5, 7, 8, 9\}$ , f(n) = 4n - 3 and g(n) = n + 2.
  - (a) Calculate f at each of 6, 10, 12, 14, 15, 30.
  - (b) Calculate f \* g at 7, 18, 30.
- 2. (Another Mersenne prime)
  - (a) Let p, q be primes such that  $q|2^p-1$ . In class we showed that  $q \equiv 1$  (p). Show that if p is odd then  $q \equiv 1 (2p)$ .

*Hint*: q - 1 is even.

- (b) Prove that  $2^{13} 1$  is prime by (i) Explicitly showing that it is not divisible by two specific primes and (ii) showing that your two trial divisions are enough.
- 3. (An amusing identity)
  - (a) Show that  $f(n) = 2^{\omega(n)}$  is a multiplicative function. *Hint*: Adapt the argument that proved that  $\mu$  was multiplicative.
  - (b) Show that  $\sum_{d|n} f(d) = \tau(n^2)$ . *Hint:* First show that it is enough to check when n is a prime power, then do that case. SUPP What would happen for  $f(d) = b^{\omega(n)}$  for a general  $b \in \mathbb{Z}_{\geq 2}$ ?
- 4. (§7.4.E30) Show that  $\Lambda * I = \log$ .

*Hint:* Use the factorization of the integer under consideration.

- 5. Define a function  $\chi_4(n) = \begin{cases} 1 & n \equiv 1 \, (4) \\ -1 & n \equiv 3 \, (4) \text{ and set } s(n) = \sum_{d|n} \chi_4(d) \text{ (i.e. } s = \chi_4 * I). \\ 0 & 2|n \end{cases}$ (a) Show that  $\chi_4$  is completely multiplicative and conclude that s is multiplicative.

  - (b) Calculate s(2), s(3), s(4), s(5).
     (c) Let r<sub>2</sub>(n) = #{(a,b) ∈ Z² | a² + b² = n} be the number of ways to write n as a sum of two squares of integers (possibly negative!). Show that  $r_2(n) = 4s(n)$  for n = 2, 3, 4, 5.

RMK The identity  $r_2(n) = 4s(n)$  holds for all n. In particular,  $r_2$  is multiplicative.

# Cryptology

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- 6. Using the affine cipher  $C \equiv 5P + 7(26)$ 
  - (a) Encrypt the message: WE ARE GOING HOME X.
  - (b) Decrypt the message: HOVYQ PBYVR VTLZZ WGZOX ZD.

7. The following message has been encoded using an affine cipher. Decode it and explain your reasoning

> NMCWT FIHHI ACPBN RSWHI NRUNG VSWBI BAUFS CPBAI YTHSI PNRSM CTSCH HYIYW UMSFS NRSTG FGAGV SWCPB NRSYG YSFCN RWGEN AFCTS

(Hint 1: the average frequency of letters in english falls according to ETAOIN)

(Hint 2: the author of passage is the Rev. C.L. Dodgson, well-known for such works as "Symbolic Logic Part I")

- 8. Show that in the following two affine ciphers encryption and decryption are the same operation (that is, that E(E(P)) = P)
  - (a) "ROT-13", a popular cipher for internet discussion boards, for which the encryption function is  $E(P) \equiv P + 13(26)$ .
  - (b) "Atbash", a historical cipher originally used in Hebrew, consisting of exchanging letters:  $a \leftrightarrow z$ ,  $b \leftrightarrow y$ ,  $c \leftrightarrow x$  and so on. Its encryption function is  $E(P) \equiv -1 - P(26)$ .
- 9. (§8.4.E6) What is the ciphertext that is produced when RSA encryption with the public key (e = 7, n = 2627) is applied to the plaintext LIFE IS A DREAM?
- 10. (§8.4.E8) In this problem your will do an RSA decryption when the public key is (e = 5, n = 2881).
  - (a) Calculate  $\phi(n)$  and find the decryption exponent d.
  - (b) If the ciphertext is 0504 1874 0347 0515 2088 2356 0736 0468, what is the plaintext message?

### **Supplementary problems (not for submission)**

- A. Fix a prime p.
  - (a) Let  $f(x) = \sum_{i=0}^{n} a_i x^i$  be a polynomial with integer coefficients. Use the identity of PS2 problem 8 to show that x - y divides f(x) - f(y) as polynomials.
  - (b) Let  $c_1 \in \mathbb{Z}$  be such that  $f(c_1) \equiv 0(p)$ . Plugging in  $c_1$  for y show that for some polynomial g(x) with integer coefficients we have a congruence of polynomials  $f(x) \equiv (x - y)^{-1}$  $c_1)g(x)(p)$ . Moreover,  $\deg(g) < \deg(f) - 1$ .
  - (c) Let  $c_2 \in \mathbb{Z}$  also be such that  $f(c_2) \equiv 0$  (p) and assume that  $c_1 \not\equiv c_2$  (p). Show that  $g(c_2) \equiv$ 0(p).
  - (d) Show by induction on r that if  $\{c_j\}_{j=1}^r$  are representatives of the distinct congruence classes mod p which solve the equation  $f(x) \equiv 0$  (p) then there is a polynomial g(x) of degree  $\leq n - r$  such that  $f(x) \equiv g(x) \prod_{i=1}^{r} (x - c_i)$ .
  - (e) Show that if f is not zero mod p then has at most n distinct roots mod p.
- B. Let f be an arithmetical function.
  - (a) Show that f is invertible (there is g such that  $f * g = \delta$ ) iff  $f(1) \neq 0$ .
  - (b) Let f be invertible. Show that it has a unique inverse and that  $(f^{-1})^{-1} = f$ . (c) Let f be invertible and multiplicative. Show that  $f^{-1}$  is multiplicative.

#### CHAPTER 4

# Cryptology

#### 4.1. Introduction

- Three parties: A (Alice) would like to communicate some message *P* ("plaintext") to B (Bob). The eavesdropper Eve will know everything Alice sends.
- Alice and Bob will agree on a pair of functions E ("encryption"), D ("decryption") such that D(E(P)) = P for all P. Alice will send the *ciphertext* C = E(P). Bob will recover the plaintext by evaluating P = D(C).
  - "Symmetric crypto": Alice and Bob keep D, E secret. Eve only knows C and has to guess what D, E are.
  - "Public-key / Asymmetric crypto": Eve knows both C and the function E, while Bob's function D is kept secret.
- The first scheme requires prior communication between Alice and Bob (to agree on the functions). Usually this prior communication is facilitated by a method of the second kind.
  - PKC is more involved, more computationally expensive. Usually only used for key exchange after which a symmetric cipher is used.
- Alice and Bob do "cryptography" (create methods of communications); Eve does "cryptanalysis" (breaking such methods).

#### 4.2. Character and block ciphers

In a *character cipher* we encode every letter of the message as an integer  $P \in \{0, ..., 25\}$  (thought of as residues mod 26). The function D, E are then maps  $\{0, ..., 25\} \rightarrow \{0, ..., 25\}$ .

EXAMPLE 115. (Caesar cipher)  $E(P) \equiv P + 3(26)$  so D(C) = C - 3(26).

• HELLO + 3 = , decrypt by -3.

EXAMPLE 116. (Affine cipher) E(P) = aP + b (26).

- Must have a invertible mod 26 for this to make sense.
- In that case  $D(C) \equiv \bar{a}(C-b) = \bar{a}C \bar{a}b$  is also an affine function
- HELLO via  $\cdot 3 7$ , decrypt via  $\cdot 9 + 63 = \cdot 9 + 11$ .

REMARK 117. (ETAOIN) Character ciphers are very weak, since they preserve the *frequency distribution* of the letters (which is highly non-uniform). They also preserve the *order* of the letters (TH most common digraph, THE most common trigraph).

Even if a different substitution is used for different plaintext letters, but the sequence of substitutions repeat it's possible to recover the block length by checking the letter distributions in residue classes.

*Block ciphers*: work with several letters are once, perhaps on a rolling basis. E.g. affine-linear map on the vector coming from several letters.

# 4.3. Asymmetric encryption: RSA

#### CHAPTER 5

#### **Primitive roots**

#### **5.1. Primitive roots**

PROBLEM 118. Solve  $x^5 \equiv 7(17)$ .

Find ord<sub>17</sub>(2) = 8 and  $6^2 \equiv 2$  so ord<sub>17</sub>(6) = 16. Now take log base 16. Similarly for  $x^5 \equiv 4$  and  $x^6 \equiv 4$ .

DEFINITION 119. For (a,m) = 1 set  $\operatorname{ord}_m(a) = \min \{ n \ge 1 \mid a^n \equiv 1 \ (m) \}$  Call r a *primitive root* mod p if  $\operatorname{ord}_p(r) = \phi(m)$ .

EXAMPLE 120. Mod 17. Mod 19.

LEMMA 121. a is a primitive root iff every invertible residue is a power of a; in this case the invertible residues are given by  $\left\{a^j\right\}_{j=0}^{\phi(m)-1}$ .

THEOREM 122. ("Discerte Logarithm") Let r be a primitive root mod m. Then the equation  $x^n \equiv r^l(m)$  has solutions iff  $(n, \phi(m))|l$  in which case there are  $(n, \phi(m))$  such solutions.

In particular, b is an nth power mod m iff  $b^{\frac{\phi(m)}{(n,\phi(m))}} \equiv 1$  (m), in which case it has  $(n,\phi(m))$  nth roots.

PROOF. Changing variables to  $x = r^t$  we need to solve  $nt \equiv l(\phi(m))$  which is a linear congruence.

THEOREM 123. There is a primitive root mod m iff m is of the form  $2, 4, p^k, 2p^k$  where p is an odd prime and  $k \ge 1$ .

#### **5.2.** Primitive roots mod p

Fix a prime p.

Two key ingredients:

PROPOSITION 124. Let  $f(x) \in \mathbb{Z}[x]$  be of degree n mod p (that is,  $p \nmid a_n$ ). Then the congruence  $f(x) \equiv 0$  (p) has at most n solutions.

PROOF. If  $f(a) \equiv 0$  (p) then x - a divides  $f \mod p$ , and continue by induction.

Proposition 125.  $\sum_{d|n} \phi(d) = n$ .

THEOREM 126. For every d|p-1 there are at d elements of order dividing d, and  $\phi(d)$  elements of order exactly d.

PROOF. There are at most d elements since x has order dividing n iff  $x^d \equiv 1$  (p). Let  $A_d = \{1 \le a \le p-1 \mid \operatorname{ord}_p(a) = d\}$ . If this set is not empty let  $a \in A_d$ . Then  $\{a^j\}_{j=0}^{d-1}$  are all distinct

so these are all elements of order dividing d. Since  $\operatorname{ord}_p(a^j) = \frac{\operatorname{ord}_p(a)}{\left(j,\operatorname{ord}_p(a)\right)}$ , exactly  $\phi(d)$  of these elements are of order d exactly. It follows that  $\#A_d \leq \phi(d)$ . Thus:

$$p-1 = \#\left\{1 \leq a \leq p-1\right\} = \#\left(\cup_{d|p-1}A_d\right) \leq \sum_{d|p-1} \#A_d \leq \sum_{d|p-1} \phi(d) = p-1 \,.$$

It follows that we must have equalities throughout, that is that  $\#A_d = \phi(d)$ . In particular for n|p-1, the elements of order dividing n are  $\bigcup_{d|n} A_d$  and there are  $\sum_{d|n} \phi(d) = n$  of them.

COROLLARY 127. There are  $\phi(\phi(p)) \ge 1$  primitive roots mod p.

# **5.3.** Primitive roots mod $p^2$ , $p^k$

Idea: linear deformation.

THEOREM 128. Let r be a primitive root mod p. Then one of r, r + p is a primitive root mod  $p^2$ .

PROOF. By Euler's Theorem,  $\operatorname{ord}_{p^2}(r)|\phi(p^2)=p(p-1)$ . Also,  $r^{\operatorname{ord}_{p^2}(r)}\equiv 1$  (p) so  $p-1=\operatorname{ord}_p(r)|\operatorname{ord}_{p^2}(r)$ . If  $\operatorname{ord}_{p^2}(r)< p(p-1)$  it must equal p-1.

Now consider  $\operatorname{ord}_{p^2}(r+pt)$ . Since  $r+pt \equiv r(p)$  the same reasoning shows that  $p-1|\operatorname{ord}_{p^2}(r+pt)|p(p-1)$ . Moreover,

$$(r+pt)^{p-1} = r^{p-1} + (p-1)r^{p-2}pt + \sum_{k=2}^{p-1} {p-1 \choose k} r^{p-1-k}t^k p^k$$

$$\equiv 1 - r^{p-2}tp(p^2).$$

Note that  $p \nmid r^{p-2}$  so as long as  $p \nmid t$  (say if t = -1) we have  $p^2 \nmid (r+pt)^{p-1} - 1$  so  $\operatorname{ord}_{p^2}(r+tp) \neq p-1$ .

THEOREM 129. [not covered in class] Let p be odd and let r be a primitive root mod  $p^2$ . Then r is a primitive root mod  $p^k$ ,  $k \ge 2$ .

PROOF. Assume by induction that  $\operatorname{ord}_{p^k} r = p^{k-1}(p-1)$ . It follows that  $r^{p^{k-2}(p-1)} \not\equiv 1$  ( $p^k$ ). Since  $r^{p^{k-2}(p-1)} \equiv 1$  ( $p^{k-1}$ ), we have  $r^{p^{k-2}(p-1)} = 1 + tp^{k-1}$  for some t not divisible by p. It follows that

$$r^{p^{k-1}(p-1)} = (1+tp^{k-1})^p = 1+tp^k + \sum_{l=2}^{p-1} \binom{p}{l} t^l p^{l(k-1)} + t^p p^{p(k-1)}.$$

Now if  $k, l \ge 2$  then  $l(k-1)+1 \ge 2(k-1)=k+1+(k-2) \ge k+1$  and  $p \mid \binom{p}{l}$  if  $2 \le l \le p-1$ . Finally, if  $p \ge 3$  then  $p(k-1) \ge 3(k-1)=k+1+2(k-2) \ge k+1$  as well so

$$r^{p^{k-1}(p-1)} \equiv 1 + tp^k \not\equiv 1 (p^{k+1}).$$

# 5.4. Discrete Log and ElGamal

See textbook.

#### Math 312: Problem Set 6 (due 14/6/11)

### **Primitive roots**

- 1. For each p find a primitive root mod p,  $p^2$ : {11, 13, 17, 19}. Justify your answers.
- 2. How many primitive roots are there mod 25? Find all of them.
- 3. (Wilson's Theorem, again)
  - (a) Let  $r = \operatorname{ord}_m(a)$  and let S be the product of the r distinct residues which are powers of a mod m. Show that  $\operatorname{ord}_m(S)$  is 1 if r is odd and 2 if r is even.
  - (b) Let p be an odd prime, and let  $k \ge 1$ . Show that the product of all invertible residues mod  $p^k$  is congruent to  $-1 \mod p^k$ .
- 4. (The quadratic character of -1) Let p be an odd prime, and let r be a primitive root mod p.
  - (a) Show that  $r^{\frac{p-1}{2}} \equiv -1$  (p), and if  $p \equiv 1$  (4) use that to find a number y such that  $y^2 \equiv -1$  (p). *Hint:* For the first part, what are the solutions to  $x^2 \equiv 1$  (p)?
  - (b) Conversely, if there is y such that  $y^2 \equiv -1(p)$  show that  $\operatorname{ord}_p(y) = 4$  and conclude that  $p \equiv 1(4)$ .
- 5. (§9.2.E12) Let p be a prime. Find the least positive residue of the product of a set of  $\phi(p-1)$  incongruent primitive roots mod p.
- 6. ElGamal
  - (a) (§10.2.E6) Using ElGamal encryption with private key (p = 2543, r = 5, a = 99), sign the message P = 2525 [use the integer k = 257] and verify the signature.
  - (b) Assume that two messages  $P_1$ ,  $P_2$  are signed using the ElGamal system with private key (p, r, a) and the same integer k with resulting signatures  $(\gamma_1, s_1)$ ,  $(\gamma_2, s_2)$ . Show that anyone observing these two signed messages can now sign a message of their choosing.

*Hint*: Consider first the case where  $s_1 - s_2$  is invertible mod p - 1.

# **Quadratic reciprocity**

- 7. Let p be an odd prime and let  $q|2^p-1$ . Recall that  $q \equiv 1$  (2p).
  - (a) We have seen before that  $\operatorname{ord}_q(2) = p$ . Use this and Euler's criterion to show that 2 is a square mod q. Conclude that  $q \equiv \pm 1(8)$ .
  - (b) Show that  $M_{17} = 2^{17} 1 < 132,000$  is prime, only trying to divide by three numbers. RMK Why is it not necessary to show that these numbers are prime?
- 8. (Math 437 Midterm, 2009)
  - (a) Let  $a \ge 3$  be odd and let  $p|a^2 2$  be prime. Show that  $p \equiv \pm 1$  (8).
  - (b) Let  $a \ge 3$  be odd. Show that *some* prime divisor of  $a^2 2$  is congruent to  $-1 \mod 8$ . *Hint*: What is the residue class of  $a^2 2 \mod 8$ ?
  - (c) Show that there are infinitely many primes congruent to  $-1 \mod 8$ .

- 9. Evaulate the following Legendre symbols.
  (a)  $\left(\frac{48}{103}\right)$ ,  $\left(\frac{3325}{14407}\right)$ ,  $\left(\frac{19382}{48397}\right)$ , using factorization and quadratic reciprocity.
  (b)  $\left(\frac{799}{37}\right)$ ,  $\left(\frac{3133}{3137}\right)$ ,  $\left(\frac{39270}{49177}\right)$ , using Jacobi symbols.
- 10. Let p be a prime such that q = 4p + 1 is also prime. Show that 2 is a primitive root mod q. *Hint:* Show that if  $\operatorname{ord}_q(2) \neq q-1$  then it must divide one of  $\frac{q-1}{2}$  and  $\frac{q-1}{p}$ , and consider those cases separately.

**Supplementary problems (not for submission)** 

#### CHAPTER 6

# **Quadratic reciprocity**

### 6.1. Quadratic residues

Fix an odd prime p.

DEFINITION 130. Let a not be divisible by p. Call a a quadratic residue mod p if there is  $x \in \mathbb{Z}$  such that  $x^2 \equiv a(p)$ . Otherwise say that x is a quadratic non-residue.

NOTATION 131. The *Legendre Symbol* is given by:

$$\left(\frac{a}{p}\right) = \begin{cases} +1 & a \text{ a quadratic residue} \\ -1 & a \text{ a quadratic nonresidue} \\ 0 & p|a \end{cases}$$

We first study  $\left(\frac{a}{p}\right)$  as a function of a.

EXAMPLE 132. List all squares mod 3, 5, 7, 11 and obtain the residues and non-residues.

PROPOSITION 133. (Euler's criterion)  $\left(\frac{a}{p}\right) \equiv a^{\frac{p-1}{2}}(p)$ .

PROOF. If p|a both sides vanish mod p. Otherwise, this is the case n=2 of Theorem 122.

EXAMPLE 134. For which primes amond 3,5,7,11 is -1 a square mod p?

COROLLARY 135. (The quadratic character of 
$$-1$$
)  $\left(\frac{-1}{p}\right) = (-1)^{\frac{p-1}{2}} = \begin{cases} +1 & p \equiv 1 \ (4) \\ -1 & p \equiv 3 \ (4) \end{cases}$ .

LEMMA 136. Let  $a, a', b \in \mathbb{Z}$  with  $a \equiv a'(p)$ . Then:

$$(1) \left(\frac{a}{p}\right) = \left(\frac{a'}{p}\right);$$

(2) 
$$\left(\frac{b^2}{p}\right) = 1$$
 if  $p \nmid b$ ;

$$(3) \left(\frac{a}{p}\right) \left(\frac{b}{p}\right) = \left(\frac{ab}{p}\right).$$

PROOF. The first two claims are true by definition  $(x^2 \equiv a(p))$  iff  $x^2 \equiv a'(p)$ . For the third, it is clear if p divides one of a, b or if at least one is a quadratic residue, but the claim that if both a, b are non-residues then ab is a residue is non-trivial. In any case using Euler's criterion it is easy to check that

$$\left(\frac{ab}{p}\right) \equiv (ab)^{\frac{p-1}{2}} \equiv a^{\frac{p-1}{2}}b^{\frac{p-1}{2}} \equiv \left(\frac{a}{p}\right) \left(\frac{b}{p}\right) (p)$$

and if signs are congruent mod p they are equal since  $p \nmid 2$ .

Another criterion:

PROPOSITION 137. (Gauss's Lemma) Let  $p \nmid a$ . Then  $\left(\frac{a}{p}\right) = (-1)^s$  where s is the number of t,  $1 \le t \le \frac{p-1}{2}$  such that the least positive residue of at is greater than  $\frac{p-1}{2}$ .

PROOF. We evalute the product  $\prod_{t=1}^{\frac{p-1}{2}}(at)=a^{\frac{p-1}{2}}\cdot\left(\frac{p-1}{2}\right)!$  in another way. For this divide the numbers  $1\leq x\leq p-1$  into pairs  $\{x,p-x\}$   $(x\neq p-x$  since p is odd). If at,at' belong to the same pair then either they are equal (at which point  $t\equiv t'(p)$ ) or opposite, at which point  $at\equiv -at'(p)$  forces  $t\equiv -t'(p)$ . Since the range  $1\leq t\leq \frac{p-1}{2}$  does not contain t,t' such that  $t\equiv -t'(p)$  and since there are exactly  $\frac{p-1}{2}$  pairs it follows that  $\prod_{t=1}^{\frac{p-1}{2}}(at)\equiv\prod_{1\leq x\leq \frac{p-1}{2}}(\pm x)(p)$  where the sign is + or - according to whether  $at\equiv x$  or  $at\equiv -x$ . By assumption we have s minus signs, so

$$a^{\frac{p-1}{2}}\left(\frac{p-1}{2}\right)! \equiv (-1)^s \prod_{x=1}^{\frac{p-1}{2}} x \equiv (-1)^s \left(\frac{p-1}{2}\right)!.$$

The factor  $\left(\frac{p-1}{2}\right)$ ! is invertible mod p and we are done by Euler's criterion.

COROLLARY 138. (The quadratic character of 2)

PROOF. Explicit count using Gauss's Lemma.

# 6.2. Quadratic reciprocity

Now consider  $\left(\frac{a}{p}\right)$  as a function of p. Euler observed that this only depends on the class of p mod 4a. Gauss eventually proved this:

THEOREM 139. (Gauss) Let p,q be distinct odd primes. Then

$$\left(\frac{p}{q}\right) \left(\frac{q}{p}\right) = (-1)^{\frac{p-1}{2}\frac{q-1}{2}} = \begin{cases} -1 & p \equiv q \equiv 3 \, (4) \\ +1 & otherwise \end{cases}.$$

PROOF. (Based on Exercise 17 to section 11.2) Let  $R = \left\{1 \le a \le \frac{pq-1}{2} \mid (a,pq) = 1\right\}$ , let  $T = \left\{q, 2q, \cdots, \frac{p-1}{2}q\right\}$  and let  $S = R \sqcup T = \left\{1 \le a \le \frac{pq-1}{2} \mid (a,p) = 1\right\}$  (if (a,p) = 1 then either (a,pq) = 1 or a is divisible by q). Finally, set  $A = \prod_{a \in R} a$ . One one hand we then have:

$$\prod_{a \in S} a = \prod_{k=0}^{\frac{q-1}{2}-1} \prod_{j=1}^{p-1} (pk+j) \cdot \prod_{j=1}^{\frac{p-1}{2}} \left( p \frac{q+1}{2} + j \right) 
\equiv (p-1)!^{\frac{q-1}{2}} \left( \frac{p-1}{2} \right)! (p) 
\equiv (-1)^{\frac{q-1}{2}} \left( \frac{p-1}{2} \right)! (p)$$

by Wilson's Theorem. On the other hand we have

$$\prod_{a \in S} a = \left(\prod_{a \in R} a\right) \cdot \left(\prod_{j=1}^{\frac{p-1}{2}} (qj)\right)$$

$$= Aq^{\frac{p-1}{2}} \left(\frac{p-1}{2}\right)!$$

$$\equiv A\left(\frac{q}{p}\right) \left(\frac{p-1}{2}\right)!(p)$$

by Euler's criterion (Proposition 133). Since  $\left(\frac{p-1}{2}\right)!$  is invertible mod p and  $\left(\frac{q}{p}\right) \equiv \pm 1 \, (p)$  we conclude

$$A \equiv (-1)^{\frac{q-1}{2}} \left(\frac{p}{q}\right) (p).$$

By symmetry we also have

$$A \equiv (-1)^{\frac{p-1}{2}} \left(\frac{q}{p}\right) (q).$$

We now evaluate  $A \mod pq$ . For this note that if x is a residue class mod pq then  $x \not\equiv -x(pq)$  since p,q are odd. It follows that for each pair  $\{x,-x\}$  of invertible residue class mod pq exactly one member belongs to A. Now let  $a \in R$  and assume that  $\bar{a} \not\equiv \pm a(pq)$ . Then exactly one of  $\bar{a}, -\bar{a}$  belongs to R and together we get either a 1 or a -1 in A. Accordingly let  $R' = \left\{1 \le a \le \frac{pq-1}{2} \mid a^2 \equiv \pm 1(pq)\right\}$ . We have shown that:

$$A \equiv \pm \prod_{q \in R'} a(pq).$$

There are 4 residues  $a \mod pq$  such that  $a^2 \equiv 1$  (p). Those are  $\pm 1$  and  $\pm u$  where  $u \equiv 1$  (p) and  $u \equiv -1$  (q). There residues contribute  $\pm u$  to A. Assume first that at least one of p,q is  $\equiv 3$  (4). Then there is not x mod that prime such that  $x^2 \equiv -1$  mod that prime; a fortiori there is no a mod pq such that  $a^2 \equiv -1$  (pq) and hence  $A \equiv \pm u(pq)$ . It follows that  $A \mod p$  and  $A \mod q$  are opposite signs. If  $p \equiv 1$  (4) and  $q \equiv 3$  (4) this means that  $-\left(\frac{p}{q}\right)$  and  $\left(\frac{q}{p}\right)$  have opposite signs (as claimed), while if  $p \equiv q \equiv 3$  (4) it means that  $-\left(\frac{p}{q}\right)$  and  $-\left(\frac{q}{p}\right)$  have opposite signs (as claimed). We are left with the case  $p \equiv q \equiv 1$  (4). Now fix  $\epsilon$ ,  $\delta$  such that  $\epsilon^2 \equiv -1$  (p) and  $\delta^2 \equiv -1$  (q). By the CRT there is p0 mod p1 such that p2 mod p3 such that p4 contains precisely one from each pair, so that p4 mod p5 and hence they are equal. It follows that  $\left(\frac{p}{q}\right)$  and  $\left(\frac{q}{p}\right)$  are congruent to the same sign mod pq, and hence they are equal.  $\square$ 

REMARK 140. Alternative proof: Let  $G = \sum_{a(p)} \left(\frac{a}{p}\right) \zeta^a$  be the Gauss sum mod p. Then  $\bar{G} = \sum_{a(p)} \left(\frac{a}{p}\right) \zeta^{-a} = \left(\frac{-1}{p}\right) G$  so  $\left(\frac{-1}{p}\right) G^2 = |G|^2 = p$  by Planchere's formula. We thus have

$$p^{\frac{q-1}{2}} = (-1)^{\frac{p-1}{2}\frac{q-1}{2}} G^{q-1}$$

and

$$Gp^{\frac{q-1}{2}} = (-1)^{\frac{p-1}{2}\frac{q-1}{2}}G^{q}$$

$$\equiv (-1)^{\frac{p-1}{2}\frac{q-1}{2}}\sum_{a(p)} \left(\frac{a}{p}\right)^{q} \zeta^{-aq}(q)$$

$$\equiv (-1)^{\frac{p-1}{2}\frac{q-1}{2}}\left(\frac{q}{p}\right)G(q).$$

Multiplying by G we find:

$$G^2\left(\frac{p}{q}\right) \equiv (-1)^{\frac{p-1}{2}\frac{q-1}{2}} \left(\frac{q}{p}\right) G^2(q).$$

Since  $G^2$  is invertible mod q we conclude

$$\left(\frac{p}{q}\right) \equiv (-1)^{\frac{p-1}{2}\frac{q-1}{2}} \left(\frac{q}{p}\right) (q)$$

as a congruence in  $\mathbb{Z}[\zeta]$ . However, if a rational number is an algebraic integer then it is an integer, so this congruence holds in  $\mathbb{Z}$  and both sides are equal.

ALGORITHM 141. *To evaluate*  $\left(\frac{a}{p}\right)$ :

- (1) Reduce a mod p to get a'.
- (2) Factor a', and eliminate any square factors.
- (3) For any prime factor q of a', relate  $\left(\frac{q}{p}\right)$  to  $\left(\frac{p}{q}\right)$  using QR, and evaluate the latter recursively.

# 6.3. The Jacobi Symbol

Better algorithm: avoid factoring.

DEFINITION 142. Let P be an odd positive integer, with prime factorization  $P = \prod_i p_i$  (the  $p_i$  need not be distinct). For  $a \in \mathbb{Z}$  the *Jacobi symbol* is the function

$$\left(\frac{a}{P}\right) \stackrel{\text{def}}{=} \prod_{i} \left(\frac{a}{p_i}\right).$$

LEMMA 143. Let  $a, a', b \in \mathbb{Z}$  with  $a \equiv a'(P)$ . Then:

$$(1) \ \left(\frac{a}{P}\right) = \left(\frac{a'}{P}\right);$$

(2) 
$$\left(\frac{b^2}{P}\right) = 1 \text{ if } (P,b) = 1;$$

(3) 
$$\left(\frac{a}{P}\right)\left(\frac{b}{P}\right) = \left(\frac{ab}{P}\right)$$
.

PROOF. These all follow from the respective properties of the Legendre symbol.

THEOREM 144. Let P,Q be odd nad positive. Then:

(1) 
$$\left(\frac{-1}{P}\right) = (-1)^{\frac{P-1}{2}};$$

(2) 
$$\left(\frac{2}{P}\right) = (-1)^{\frac{P^2-1}{8}};$$

(3) 
$$\left(\frac{Q}{P}\right)\left(\frac{P}{Q}\right) = (-1)^{\frac{P-1}{2}\frac{Q-1}{2}}.$$

PROOF. In both (1), (2) the claim holds for P prime. Both sides of the claimed equality are also completely multiplicative (clear on the Jacobi Symbol side and an easy calculation on the other) so the claim follows). For part (3) one checks that both sides are separately completely multiplicative in P, Q (for the RHS this is already checked for part (1)) so again equality follows from the case of primes, which is the law of QR. 

ALGORITHM 145. To evaluate  $\left(\frac{a}{P}\right)$ :

- (1) Reduce a mod P to get a'.
- (2) Write  $a' = 2^t Q$  with Q odd. We then have  $\left(\frac{a}{P}\right) = \left(\frac{a'}{P}\right) = \left(\frac{2}{P}\right)^t \left(\frac{Q}{P}\right)$ . (3) Evalute  $\left(\frac{2}{P}\right)$  by part (2), and  $\left(\frac{Q}{P}\right)$  by relating it to  $\left(\frac{P}{Q}\right)$  and continuing recursively.

#### CHAPTER 7

# **Special topics**

### 7.1. The Gaussian integers

Baseically a review of the course but using a different number system.

- Defined the rings  $\mathbb{Q}(i)$  and  $\mathbb{Z}[i]$ .
- State that they are rings.
- Define conugation and the norm; relate it to divisibility in  $\mathbb{Q}(i)$ .
- Find all units of  $\mathbb{Z}[i]$ .
- Divisition with remainder in  $\mathbb{Z}[i]$  by rounding quotient in  $\mathbb{Q}(i)$ .
- Divisibility:
  - Only finitely many divisors;
  - define *gcd*;
  - Euclid's Lemma still holds. Due to division with remainder Euclid's Algorithm also still holds.
  - Bezout's extension also holds, and can also proof Bezout's Theorem by considering
    a minimal element of the ideal generated by z, w. This also shows the GCD is unique
    up to associates.
- Unique factorization
  - Define *irreducible*, *prime*. Discuss associates.
    - \* If Nz = p is a rational prime then z must be irreducible.
  - Show that every Guassian integer is a product of irreducibles.
  - Show that  $\pi$  is prime iff it is irreducible and not a unique.
  - Conclude that the prime factorization is unique up to permutation and associates.
- Classification of primes.
  - If  $\pi$  is prime then  $\pi$  divides  $N\pi$  which is a rational integer, and hence a product of rational primes. It follows that  $\pi|p$  for some rational prime p.
  - If  $N\pi = p^2 = Np$  then  $\pi$  is assoc to p. Otherwise  $N\pi = p$ .
  - $-2 = -i(1+i)^2$  so 1+i is the only prime dividing 2.
  - If  $p \equiv 3$  (4) and  $p|a^2 + b^2$  then p|a, p|b (if  $p \nmid a$  then  $(\bar{a}b)^2 \equiv -1$  (p)). It follows that p is not a norm, so p is still prime in  $\mathbb{Z}[i]$ .
  - If  $p \equiv 1$  (4) then there is  $a \in \mathbb{Z}$  such that  $p \mid a^2 + 1$ . But  $p \nmid a \pm i$  in  $\mathbb{Z}[i]$  so p is not prime. It follows that there is a prime properly dividing p so  $N\pi = p$ . This says  $p = \pi \bar{\pi}$  and the two are not associates:  $\frac{\pi}{\bar{\pi}} = \frac{\pi^2}{p} \notin \mathbb{Z}[i]$ , so p is divisible by exactly two primes (and we can write  $p = a^2 + b^2$  in 8 ways, coming from the 4 associates of  $\pi, \bar{\pi}$  each).

# 7.2. Elliptic curves

$$y^2 = x^3 + ax + b.$$

- Plane cubics; the addition law.
  Fermat descent for x<sup>4</sup> + y<sup>4</sup> = z<sup>2</sup>.
  Modularity
- Elliptic curves mod p.
- Elliptic curve cryptography.

# **Bibliography**