

# Forbidden Configurations: Boundary Cases

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Consider the following family of subsets of  $\{1, 2, 3, 4\}$ :

$$\mathcal{A} = \{\emptyset, \{1, 2, 4\}, \{1, 4\}, \{1, 2\}, \{1, 2, 3\}, \{1, 3\}\}$$

The incidence matrix  $A$  of the family  $\mathcal{A}$  of subsets of  $\{1, 2, 3, 4\}$  is:

$$A = \begin{bmatrix} 0 & 1 & 1 & 1 & 1 & 1 \\ 0 & 1 & 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 \\ 0 & 1 & 1 & 0 & 0 & 0 \end{bmatrix}$$

**Definition** We say that a matrix  $A$  is *simple* if it is a  $(0,1)$ -matrix with no repeated columns.

**Definition** We define  $\|A\|$  to be the number of columns in  $A$ .

$$\|A\| = 6 = |\mathcal{A}|$$

**Definition** Given a matrix  $F$ , we say that  $A$  has  $F$  as a *configuration* if there is a submatrix of  $A$  which is a row and column permutation of  $F$ .

$$F = \begin{bmatrix} 0 & 0 & 1 & 1 \\ 0 & 1 & 0 & 1 \end{bmatrix} \in A = \begin{bmatrix} 0 & 1 & 1 & 1 & 1 & 1 \\ 0 & \boxed{1} & \boxed{0} & \boxed{1} & 1 & \boxed{0} \\ 0 & 0 & 0 & 0 & 1 & 1 \\ 0 & \boxed{1} & \boxed{1} & \boxed{0} & 0 & \boxed{0} \end{bmatrix}$$

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We consider the property of forbidding a configuration  $F$  in  $A$ .

**Definition** Let

$$\text{forb}(m, F) = \max\{\|A\| : A \text{ } m\text{-rowed simple, no configuration } F\}$$

# A Product Construction

The building blocks of our product constructions are  $I$ ,  $I^c$  and  $T$ , e.g:

$$I_4 = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}, \quad I_4^c = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix}, \quad T_4 = \begin{bmatrix} 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

# The Conjecture

We conjecture that our product constructions with the three building blocks  $\{I, I^c, T\}$  determine the asymptotically best constructions.

**Definition** Given two matrices  $A, B$ , we define the product  $A \times B$  as the matrix whose columns are obtained by placing a column of  $A$  on top of a column of  $B$  in all possible ways. (A, Griggs, Sali 97)

$$\begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} \times \begin{bmatrix} 1 & 1 & 1 \\ 0 & 1 & 1 \\ 0 & 0 & 1 \end{bmatrix} = \left[ \begin{array}{ccc|ccc|ccc} 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 \\ \hline 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 0 & 1 & 1 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 & 0 & 1 & 0 & 0 & 1 \end{array} \right]$$

Given  $p$  simple matrices  $A_1, A_2, \dots, A_p$ , each of size  $m/p \times m/p$ , the  $p$ -fold product  $A_1 \times A_2 \times \dots \times A_p$  is a simple matrix of size  $m \times (m/p)^p$  i.e.  $\Theta(m^p)$  columns.

# Examples

$$[01] \times [01] = K_2 = \begin{bmatrix} 1 & 1 & 0 & 0 \\ 1 & 0 & 1 & 0 \end{bmatrix}$$

$I_{m/2} \times I_{m/2}$  is vertex-edge incidence matrix of  $K_{m/2,m/2}$



# The Conjecture

We conjecture that our product constructions with the three building blocks  $\{I, I^c, T\}$  determine the asymptotically best constructions.

**Definition** Let  $F$  be given. Let  $x(F)$  denote the largest  $p$  such that there is a  $p$ -fold product which does not contain  $F$  as a configuration where the  $p$ -fold product is  $A_1 \times A_2 \times \cdots \times A_p$  where each  $A_i \in \{I_{m/p}, I_{m/p}^c, T_{m/p}\}$ .

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The conjecture has been verified for  $k \times \ell$   $F$  where  $k = 2$  (A, Griggs, Sali 97) and  $k = 3$  (A, Sali 05) and  $\ell = 2$  (A, Keevash 06) and for  $k$ -rowed  $F$  with bounds  $\Theta(m^{k-1})$  or  $\Theta(m^k)$  (A, Fleming 10) plus other cases.

**Definition** Let  $F$  be a  $k$ -rowed configuration and let  $\alpha$  be a  $k$ -rowed column vector. Define  $[F|\alpha]$  to be the concatenation of  $F$  and  $\alpha$ .

**Definition** Let  $F$  be a  $k$ -rowed configuration. We say that  $F$  is a **boundary case** if for every  $k$ -rowed column  $\alpha$  which is either  
not present in  $F$

or

present just once in  $F$ ,

then  $forb(m, [F|\alpha])$  is  $\Omega(m \cdot forb(m, F))$ .

# Example of a Boundary Case

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$$F = \begin{bmatrix} 0 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 \\ 0 & 1 & 1 & 1 \end{bmatrix}, \quad \text{forb}(m, F) = 2m$$

Thus  $F$  is a **boundary case**.

Using a result of A and Fleming 10, there are three simple **column-maximal** 4-rowed  $F$  for which  $\text{forb}(m, F)$  is quadratic. Here is one example:

$$F_8 = \begin{bmatrix} 1 & 0 & 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 & 0 & 0 \end{bmatrix}$$

How can we repeat columns in  $F_8$  and still have a quadratic bound? We note that repeating either the column of sum 1 or the column of sum 3 will result in a cubic lower bound. Thus we only consider taking multiple copies of the columns of sum 2.

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How can we repeat columns in  $F_8$  and still have a quadratic bound? We note that repeating either the column of sum 1 or the column of sum 3 will result in a cubic lower bound. Thus we only consider taking multiple copies of the columns of sum 2. For a fixed  $t$ , let

$$F_8(t) = \left[ \begin{array}{cccc} 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 1 \end{array} t \cdot \begin{array}{cc} 1 & 0 \\ 0 & 1 \\ 1 & 1 \\ 0 & 0 \end{array} \right]$$

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**Theorem** (A, Raggi, Sali 09) *Let  $t$  be given. Then  $\text{forb}(m, F_8(t))$  is  $\Theta(m^2)$ . Moreover  $F_8(t)$  is a **boundary case**, namely for any column  $\alpha$  not already present  $t$  times in  $F_8(t)$ , then  $\text{forb}(m, [F_8(t)|\alpha])$  is  $\Omega(m^3)$ .*

The proof of the upper bound is currently a rather complicated induction with some directed graph arguments.

## 5 × 6 Simple Configuration with Quadratic bound

The Conjecture predicts nine 5-rowed simple matrices  $F$  to be **boundary cases**, namely  $\text{forb}(m, F)$  is predicted to be  $\Theta(m^2)$  and for any column  $\alpha$  we have  $\text{forb}(m, [F|\alpha])$  being  $\Omega(m^3)$ . We have handled the following case.

$$F_7 = \begin{bmatrix} 1 & 1 & 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 1 & 1 & 1 \\ 0 & 1 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 & 0 \end{bmatrix}$$

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# All 6-rowed Configurations with Quadratic Bounds

$$G_{6 \times 3} = \begin{bmatrix} 1 & 1 & 1 \\ 1 & 1 & 0 \\ 1 & 0 & 1 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \\ 0 & 0 & 0 \end{bmatrix}$$

**Theorem** (A,Raggi,Sali)  $\text{forb}(m, G_{6 \times 3})$  is  $\Theta(m^2)$ . Moreover  $G_{6 \times 3}$  is a *boundary case*, namely for any column  $\alpha$ , then  $\text{forb}(m, [G_{6 \times 3} | \alpha])$  is  $\Omega(m^3)$ . In fact if  $F$  is not a configuration in  $G_{6 \times 3}$ , then  $\text{forb}(m, F)$  is  $\Omega(m^3)$ .

**Proof:** We use induction and the bound for  $F_7$ .

**Theorem** (Balogh and Bollabás 05) *Given  $k$ , there exists a constant  $c_k$  so that  $\text{forb}(m, \{I_k, I_k^c, T_k\}) = c_k$ .*

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**Theorem** (A. and Meehan 11) *Let  $p, k$  be given with  $p \geq 3k$ . Let  $F = [0_k | I_k] \times [0_k | T_k] \times [I_k^c | 1_k] \times K_{p-3k}$ . Then  $\text{forb}(m, F)$  is  $\Theta(m^{p-k})$ .*

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*Then  $\text{forb}(m, F)$  is  $\Theta(m^{p-k})$ .*

*Moreover  $F$  is **column maximal** (a **weak form** of a **boundary case**), namely for any column  $\alpha$  not in  $F$  we have that  $\text{forb}(m, [F | \alpha])$  is  $\Omega(m^{p-k+1})$ .*

# Induction

Let  $A$  be an  $m \times \text{forb}(m, F_7)$  simple matrix with no configuration  $F_7$ . We can select a row  $r$  and reorder rows and columns to obtain

$$A = \text{row } r \begin{bmatrix} 0 & \cdots & 0 & 1 & \cdots & 1 \\ B_r & & C_r & C_r & & D_r \end{bmatrix}.$$

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Now  $[B_r C_r D_r]$  is an  $(m - 1)$ -rowed simple matrix with no configuration  $F_7$ . Also  $C_r$  is an  $(m - 1)$ -rowed simple matrix with no configurations in  $\mathcal{F}$  where  $\mathcal{F}$  is derived from  $F_7$ .

$C_r$  has no  $F$  in

$$\mathcal{F} = \left\{ \begin{array}{l} \begin{bmatrix} 1 & 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 \end{bmatrix}, \begin{bmatrix} 1 & 1 & 0 & 1 \\ 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}, \begin{bmatrix} 1 & 1 & 0 & 1 & 0 \\ 1 & 0 & 1 & 1 & 1 \\ 0 & 1 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 & 1 \end{bmatrix} \end{array} \right\}$$

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$$\|A\| = \text{forb}(m, F_7) = \|B_r C_r D_r\| + \|C_r\| \leq \text{forb}(m-1, F_7) + \|C_r\|.$$

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To show  $\text{forb}(m, F_7)$  is quadratic it would suffice to show  $\|C_r\|$  is linear for some choice of  $r$ .

# Repeated Induction

Let  $C_r$  be an  $(m - 1)$ -rowed simple matrix with no configuration in  $\mathcal{F}$ . We can select a row  $s_j$  and reorder rows and columns to obtain

$$C_r = \text{row } s_j \begin{bmatrix} 0 & \cdots & 0 & 1 & \cdots & 1 \\ E_j & & G_j & G_j & & H_j \end{bmatrix}.$$

To show  $\|C_r\|$  is linear it would suffice to show  $\|G_j\|$  is bounded by a constant for some choice of  $s_j$ . Our proof shows that assuming  $\|G_j\| \geq 8$  for all choices  $s_j$  results in a contradiction.

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Then we discover:

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We may choose  $s_1$  and form  $L_1$ .

Then choose  $s_2 \in L_1$  and form  $L_2$ .

Then choose  $s_3 \in L_2$  and form  $L_3$ .

etc.

We can show the sets  $L_1 \setminus s_2, L_2 \setminus s_3, L_3 \setminus s_4, \dots$  are disjoint.

Assuming  $\|G_i\| \geq 8$  for all choices  $s_i$  results in  $|L_i \setminus s_{i+1}| \geq 3$  which yields a contradiction.

THANKS to the organizers, particularly Sali Attila!