# Forbidden Configurations 

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- At most how many $2 \times 1$ dominoes can we place on a chessboard which has two opposite corners removed?
Answer: 30
- At most how many edges can a simple graph with $p$ vertices have, if it has no triangles?
Answer: $\left\lfloor\frac{p^{2}}{4}\right\rfloor$ (Turán's Theorem)

Definition. A simple matrix is a $\{0,1\}$-matrix with no repeated columns.

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$$
\text { e.g. }\left[\begin{array}{llllll}
0 & 1 & 1 & 0 & 0 & 0 \\
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0 & 0 & 1 & 1 & 1 & 1
\end{array}\right]
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An m-rowed simple matrix has at most $2^{m}$ columns.

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Definition. Suppose that $F$ is a $\{0,1\}$-matrix (not necessarily simple). A simple matrix $A$ has the configuration $F$ if $A$ has a submatrix which is a row and column permutation of $F$.

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\end{array}\right]
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Extremal Problem: If a simple matrix $A$ has $m$ rows and does not have the configuration $F$, at most how many columns can $A$ have?

Answer: forb $(m, F)$

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Definition. Suppose that $F$ is a $\{0,1\}$-matrix, and $m$ a positive integer. Then forb $(m, F)$ is the greatest number of columns that an $m$-rowed simple matrix with no configuration $F$ can have.

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Equivalently, forb $(m, F)$ is the least integer such that every simple matrix with $m$ rows and more than forb $(m, F)$ columns has the configuration $F$.

## Examples

$$
\text { forb }\left(m,\left[\begin{array}{ll}
1 & 0
\end{array}\right]\right)=1 .
$$

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$\operatorname{forb}\left(m,\left[\begin{array}{ll}1 & 0\end{array}\right]\right)=1$.
forb $\left(m,\left[\begin{array}{l}1 \\ 1 \\ 0 \\ 0\end{array}\right]\right)=2 m+2$.

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\begin{aligned}
& \operatorname{forb}\left(m,\left[\begin{array}{ll}
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\end{array}\right]\right)=1 \\
& \text { forb }\left(m,\left[\begin{array}{l}
1 \\
1 \\
0 \\
0
\end{array}\right]\right)=2 m+2 \\
& \text { forb }\left(m,\left[\begin{array}{ll}
1 & 0 \\
0 & 1
\end{array}\right]\right)=m+1
\end{aligned}
$$

## Some New Results for 2-Columned F

Others proved previously that

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\end{array}\right]\right)=m+1 \text {, }
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\text { forb }\left(m,\left[\begin{array}{ll}
1 & 0 \\
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\end{array}\right]\right)=\binom{m}{2}+m+2 \quad \forall m \geq 3 .
\end{gathered}
$$

What happens if we keep adding $\left[\begin{array}{ll}1 & 0\end{array}\right]$ on top?

Theorem. For $m \geq 3$,

$$
\text { forb }\left(m,\left[\begin{array}{ll}
1 & 0 \\
1 & 0 \\
1 & 0 \\
0 & 1
\end{array}\right]\right)=\binom{m}{2}+m+2
$$

A construction: the m-rowed matrix with all columns of sum

$$
0,1,2 \text { and } m .
$$

Theorem. For $m \geq 4$,

$$
\text { forb }\left(m,\left[\begin{array}{cc}
1 & 0 \\
1 & 0 \\
1 & 0 \\
1 & 0 \\
0 & 1
\end{array}\right]\right)=\binom{m}{3}+\binom{m}{2}+m+2
$$

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0,1,2,3 \text { and } m .
$$

Theorem. For $m \geq 5$,

$$
\text { forb }\left(m,\left[\begin{array}{cc}
1 & 0 \\
1 & 0 \\
1 & 0 \\
1 & 0 \\
1 & 0 \\
0 & 1
\end{array}\right]\right)=\binom{m}{4}+\binom{m}{3}+\binom{m}{2}+m+2
$$

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0,1,2,3,4 \text { and } m .
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## Theorem. For $m \geq k-1 \geq 3$,



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I asked, "What if I flip some digits in the second column?"
Theorem. For $m \geq k-1 \geq 3$,

$$
\text { forb } \left.\left(m,\left[\begin{array}{cc}
1 & 0 \\
1 & 0 \\
1 & 0 \\
\vdots & \vdots \\
1 & 0 \\
1 & 0 \\
1 & 0 \\
0 & 1
\end{array}\right]\right\} k\right)=\binom{m}{k-2}+\cdots+\binom{m}{2}+m+2 .
$$

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0,1,2, \ldots, k-2 \text { and } m .
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The bound and the construction remains the same!
New! Theorem. For $m \geq k-1 \geq 3$,

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Finding a good construction becomes a difficult Design Theory problem.

What if we flip the 1 at the bottom of the second column to a 0 ?
For $m \geq k-1 \geq 3$,


## We get the same result as before.

New! Theorem. For $m \geq k-1 \geq 3$,

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\text { forb } \left.\left(m,\left[\begin{array}{cc}
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## Sketch of Induction Proof:

For a given $F$, let $A$ be a simple $m \times$ forb $(m, F)$ matrix which does not have the configuration $F$.

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We permute the columns of $A$ so that it looks like

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\left[\begin{array}{ccccc}
0 & \cdots & 0 & 0 & 1
\end{array} 1 \cdots \cdots 111 .\right.
$$

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11 & 1 & \cdots & 1 \\
C & D & D & E
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Then $C, D, E$ concatenated together is simple and $(m-1)$-rowed, and does not have the configuration $F$.

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\begin{aligned}
\therefore & \# \mathrm{col} \text { 's }(A)=\# \operatorname{col} \text { 's }(C, D, E)+\# \operatorname{col} \text { 's }(D) \\
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If we can get a good upper bound on $\#$ col's $(D)$, then we can prove an upper bound on forb $(m, F)$ by induction.

## Thank You!

Thanks for your attention!

