

Local & global structure of uniform spanning trees

(joint works with subsets of {Noga Alon, Eleanor Archer,
Peleg Michaeli, Yuval Paoes,
Matan Shalev})

OOPS 2021

Setup: G finite, connected, simple graph

T is a spanning tree if $V(T) = V(G)$

T conn.

no cycles.

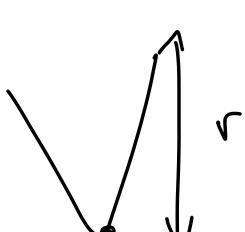
Rick: "The UST is a hard
prob. theory"

(1859)

Cayley's theorem : # trees on n labelled vertices
 $= n^{n-2}$

Kirchhoff's thm (matrix-tree)

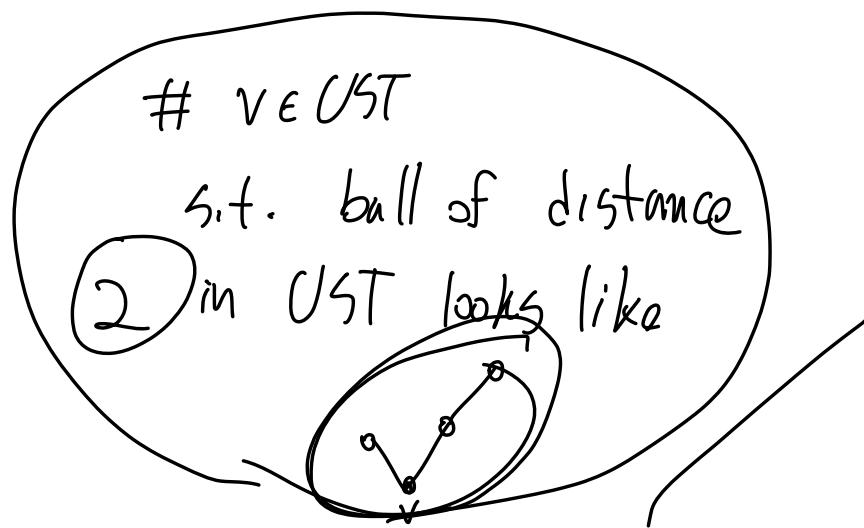
$$\# \text{ spanning trees of any } G = \det \begin{pmatrix} \text{minor} \\ \text{the Laplacian} \end{pmatrix}$$

- Global structure :
- (+) Typical dist between a pair of vertices
 - (+) Diameter (len longest path)
 - (+) Height seen from random vertices
 - / (+) The Does the \parallel
- 

tree, considered as
a metric space, converge.

(Gromov-Hausdorff-Prokhorov)

local structures $\left. \begin{array}{l} \oplus \# \text{ leaves} \\ \oplus \# \text{ deg } k \text{ vertices} \end{array} \right\}$



\oplus local limit of the UST

Case study: UST(K_n)

$T = UST(K_n)$

$\gamma^2,$

Claim: $x \neq y \quad P\left(d_T(x, y) \geq \sqrt{n}\right) = e^{-n/2} + o(1)$

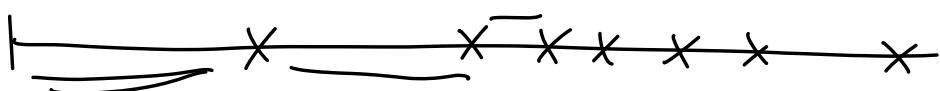
tree distance

Renyi - Szekeres (67) $\frac{\text{Height}}{\sqrt{n}} \xrightarrow{(d)} \text{cont r.v. supp on } \mathbb{S}$

Szekeres (82) : $\frac{\text{diameter}}{\sqrt{n}} \xrightarrow{(d)}$

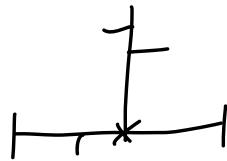
Ultimate theorem goes (Abdus 90-93)
Le Gall

$\frac{d_T(\cdot, \cdot)}{\sqrt{n}} \xrightarrow[\text{GHP}]{(d)} \text{Continuum Random Tree (CRT)}$



\mathbb{R}^+ Poisson process w/ intensity $t dt$
 $\cdot L_{t,x,t}$

(line breaking)



Brownian excursion

Corollary: $\frac{\text{Height}}{\sqrt{n}} \xrightarrow{(d)} \sqrt{2} \cdot \max_{t \in [0,1]} l_t$

$\frac{\text{Drum}}{\sqrt{n}} \xrightarrow{(d)} \sqrt{2} \sup_{\begin{array}{l} 0 \leq t_1 < t_2 \leq 1 \\ t_2 \leq t \leq t_1 \end{array}} l_{t_2} + l_{t_1} - 2 \inf_{t_2 \leq t \leq t_1} l_t$

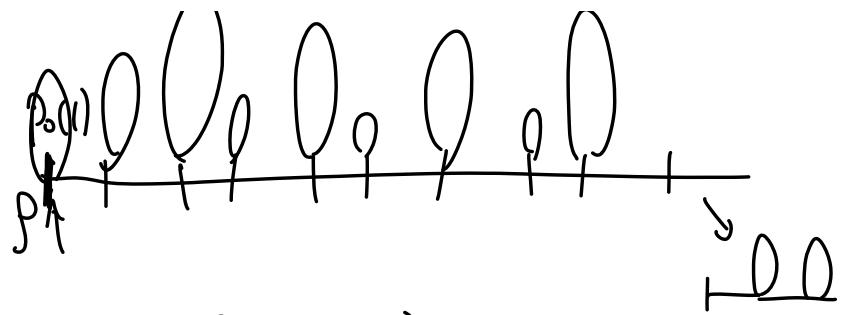
Local structure $UST(k_n) = T$

Prufer codes degree $\stackrel{(d)}{=} \text{Bin}(n-2, \frac{1}{n}) + 1$

Ultimate theorem (Grimmett)
81

Local limit of T is the $po(1)$

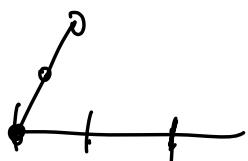
BP conditioned to survive.



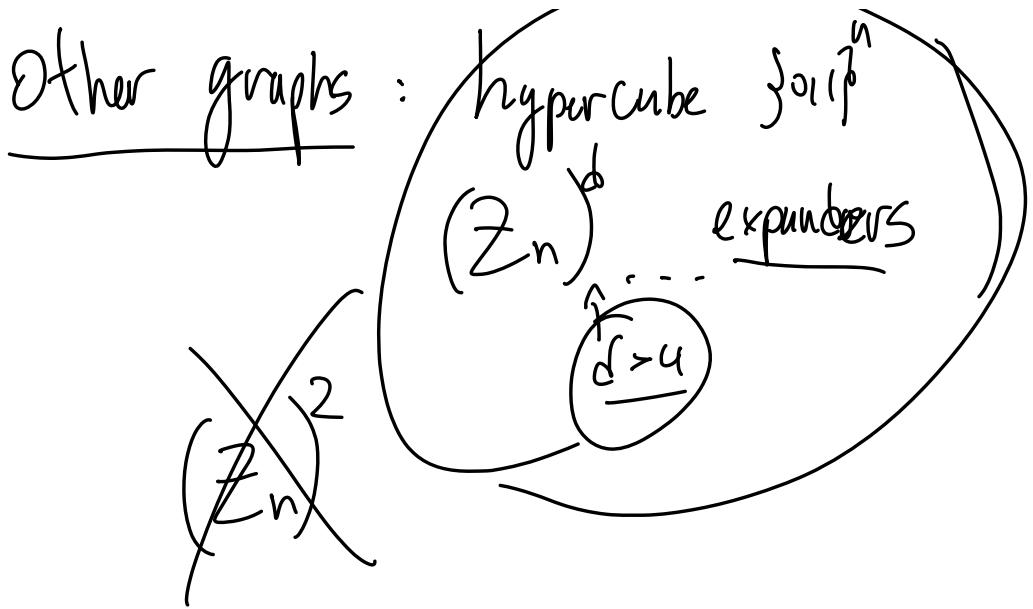
Corollary : $P(\text{v leaf}) = \frac{1}{e}$

$$\begin{aligned} \# \text{ leaves} &= \frac{n}{e}(1+o(1)) \quad \text{whp} \\ \deg \text{ root} &\stackrel{(d)}{=} 1 + P_0(1) \\ \# \text{ } \times \text{ } \times &= \frac{n}{e^3}(1+o(1)) \end{aligned}$$

$$\# \text{ } \text{ } \text{ } \text{ } = \frac{n}{e^3}(1+o(1))$$



All this was for K_n



Results Global

Thm (Michael, N., Shalev)

For any "high dimensional" graph \mathbb{G} ,

$$\forall \epsilon \exists A \quad P\left(\underbrace{A\sqrt{n}}_{\text{UST}(G_n)} \leq \text{diam}(\mathbb{G}) \leq A\sqrt{n}\right) \geq 1 - \epsilon$$

$$n = \# V(\mathbb{G}_n)$$

Work in progress (Archer, N., Shalev)

Same setup, $\exists \beta > 0$ s.t.

$$\frac{\text{diam}(T_n)}{\sqrt{n}} \xrightarrow{(a)} \text{diam}(\text{CRT})$$

$G = (Z_n)^S$

in fact $\beta \cdot \frac{d_{T_n}(\cdot, \cdot)}{\sqrt{n}} \xrightarrow{(b)} \text{CRT}$

Thm (Peres, N.) G_n a graph seq., regular
w.r.t. $\deg \rightarrow \infty$

then local limit of $\text{UST}(G_n)$

is $\text{Po}(1)$ BP cond. survival

Thm (Kirchhof 1847): \mathbb{F} finite conn. graph
 $e = (x, y) \in E(\mathbb{G})$, $T = UST(\mathbb{F})$

$$P(e \in T) = \frac{R_{\text{eff}}(x \leftrightarrow y)}{\dots}$$

Pf: $\theta(\vec{e}) := P\left(\begin{array}{l} \text{the unique } x \rightarrow y \\ \text{path in } T \\ \text{uses } \vec{e} \end{array}\right) - P\left(\begin{array}{l} " \\ \vec{e} \end{array}\right)$

Claim: θ is $\frac{\text{unit current}}{\sqrt{V}} \rightsquigarrow \frac{\text{flow}}{\sqrt{V}}$.

Why cycle law holds, i.e., $\nabla \vec{e}_1, \dots, \vec{e}_m$ cycle

$$\sum_{i=1}^m \theta(\vec{e}_i) \stackrel{?}{=} 0$$

t spanning tree of \mathbb{G}

$$\sum_t \sum_{i=1}^m f_i^+(t) - \sum_t \sum_{i=1}^m f_i^-(t)$$

\dots if \vec{n} is island

$$f_i^+(t) = \begin{cases} +1 & \text{if } x \rightarrow y \text{ in path in } t \\ 0 & \text{otherwise} \end{cases}$$

$$f_i^-(t)$$

$$T_i^+ = \left\{ (t, i) : t \text{ span-tree}, f_i^+(t) = 1 \right\}$$

$$T_i^- = \quad "$$

$$\left| \bigcup_{i=1}^m T_i^+ \right| = \left| \bigcup_{j=1}^m T_j^- \right|$$



$t \setminus \{e_i\}$ has 2 conn. comp

Let e_j be the 1st edge after \vec{e}_i (in the cycle)
that connects these 2 comp.

$$\Rightarrow \underline{(t \setminus \{e_i\} \cup \{e_j\}, j)} \in T_j^-$$

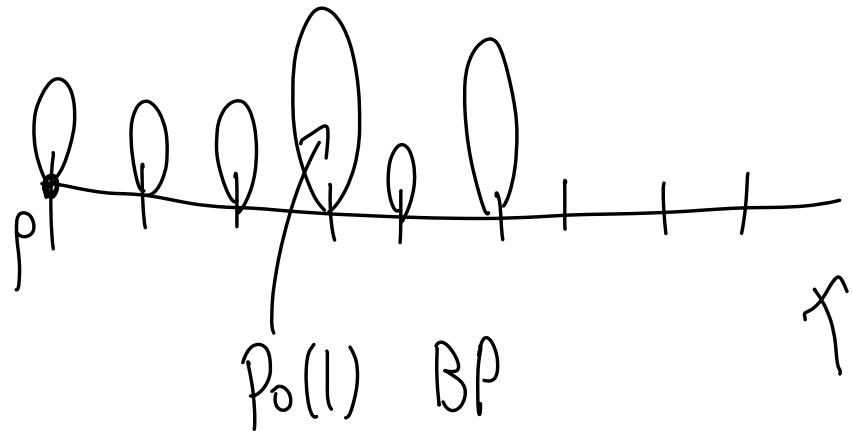
□

The local structure of USTs

Thm (Yuval Peres, N.)
+ 21

G_n sequence of connected, finite
simple, regular with $\deg \rightarrow \infty$.

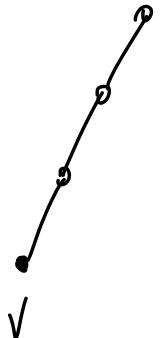
Then $\underline{UST}(G_n)$ converges locally
Po(l) BP conditions to survive.



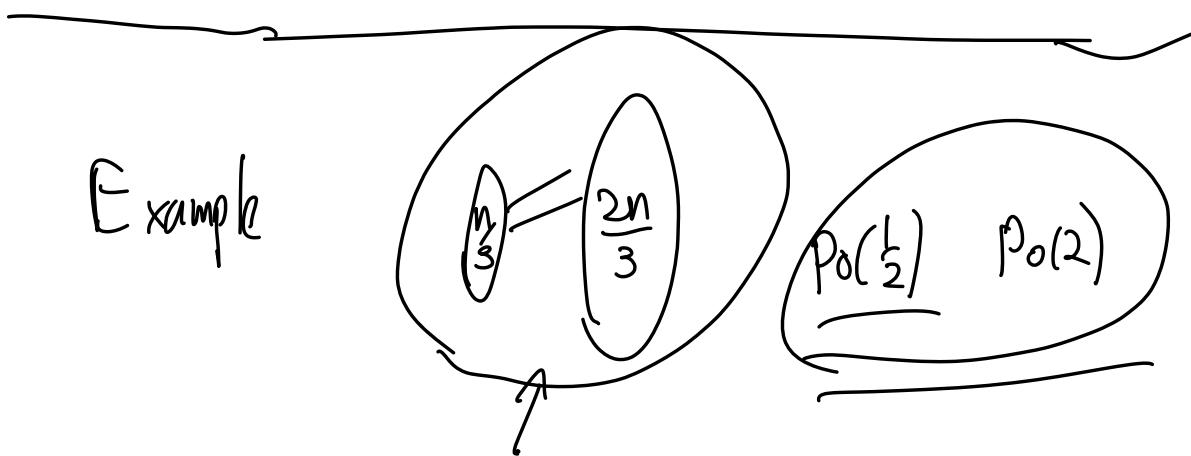
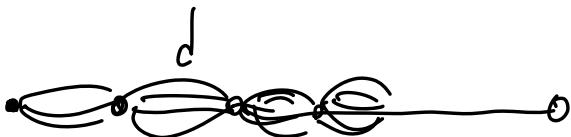
I.e. x uniform vertex of G_n

Then $\lim_{n \rightarrow \infty} P\left(\frac{B(X_1, r)}{T_n} = T\right) = P\left(\frac{B(X_1, r) = T}{P(1) / P(T)}\right)$

$$T_n = \text{UST}(F_n)$$



$$\frac{n}{c^3} (1 + o(1))$$



Reminder about electric networks

$c(i,j)$ = conductance on (i,j)

$\bar{c}(i) = \sum_j c(i,j)$

$$\textcircled{+} \quad R_{\text{eff}}(a \leftrightarrow z) = \frac{1}{T_a P_a(\tau_z < \tau_a^+)} \quad |$$

$\textcircled{+}$ Commute time identity (in unit cond)

$$2E(G) R_{\text{eff}}(a \leftrightarrow z) = \underline{\underline{E_a \tau_z + E_z \tau_a}}$$

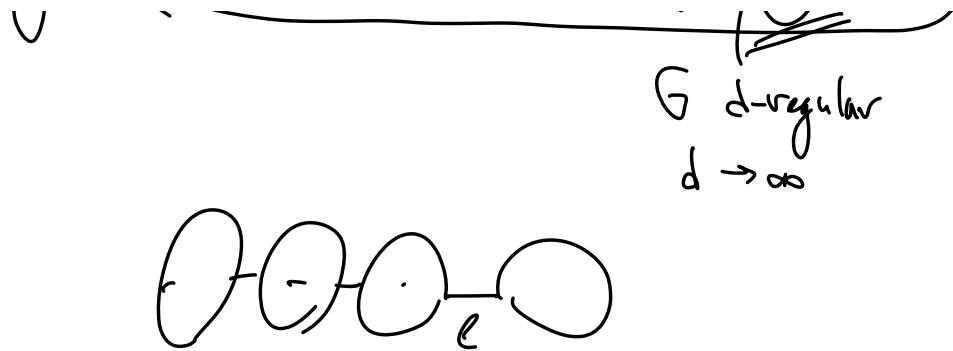
Kirchhoff (1847) : $e \in E \quad e = (x, y)$

$$P(e \in T) = R_{\text{eff}}(x \leftrightarrow y)$$

Foster's theorem: $\sum_{e \in (x, y)} R_{\text{eff}}(x, y) = n - 1$

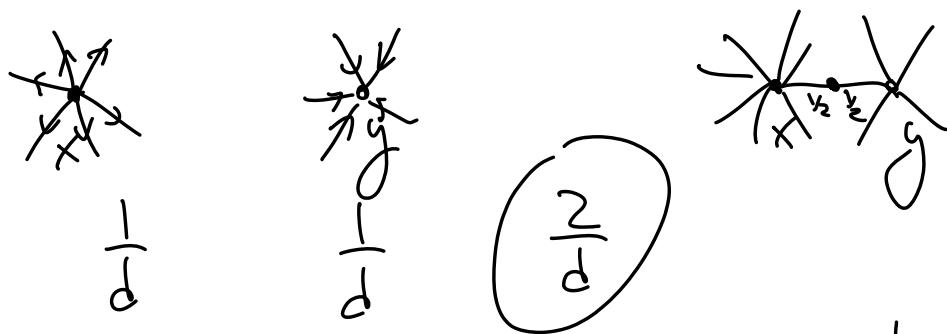
immediate

Alternatively: $\boxed{E R_{\text{eff}}(x_0, x_1) = \frac{n-1}{E(G)} = \left(\frac{2}{d}\right)(1 + \alpha_1)}$



Observation: In any graph, $x \neq y$

$$R_{\text{off}}(x \rightarrow y) \geq \frac{1}{\deg(x)+1} + \frac{1}{\deg(y)+1}$$



Corollary $N_\varepsilon = \left| \{e = (x,y) : R(x,y) \geq \varepsilon \} \right|$

$$\left| N_\varepsilon \right| \leq \frac{n}{ed - 2} \quad \varepsilon > \frac{2}{d}$$

A random edge $\underline{(X_0, X_1)}$ whp $R(X_0, X_1) = \frac{2}{d}(1+\alpha)$

X_0 - random vertex X_1 - random nbr

(X_0, \dots, X_k) - SRW k -steps

Goal: whp $R(X_0, X_k) = \frac{2}{d}(1+\alpha)$

Proof of

$$\text{Fix } X_0 \in V \quad \mathbb{E}_{X_0} \tau_{X_0^+} = n = \underbrace{\frac{1}{d^k} \sum_{\substack{x_0, x_1, \dots, x_k \\ x_i \sim x_0}} \left(\sum_i (\mathbb{E}_{X_i} \tau_{X_0^+} + 1) \right)}_{\text{oval}}$$

$$n-1 = \frac{1}{d} \sum_{X_1 \sim X_0} \mathbb{E}_{X_0} \tau_{X_1} = \mathbb{E} \mathbb{E}_{X_1} \tau_{X_0}$$

↑
random nbr
of X_0

X_0 - uniform vertex

$$\mathbb{E} \sum_{\substack{i=1 \\ i \neq k}}^n \mathbb{P}_{X_0, X_k} + \mathbb{E}_{X_0, X_k} \mathbb{P}_{X_0, X_k} = 2(n-1) = 2\mathbb{E}(G)R_{\text{eff}} \quad (X_0 \leftrightarrow X_k)$$

Similarly $\mathbb{E} R_{\text{eff}}(X_0 \leftrightarrow X_k) = \frac{2}{d} (1+o(1))$

Remark: again whp on the RW

$$R_{\text{eff}}(X_0, X_k) = \left(\frac{2}{d}\right)(1+o(1))$$

Thm

$$\mathbb{X}_1, \dots, \mathbb{X}_k \quad k \geq 2$$

whp

$$R_{\text{eff}}(X_k \leftrightarrow \{X_1, \dots, X_{k-1}\}) = \frac{k}{(k-1)d} (1+o(1))$$



We know $\forall i \neq j \in [k]$ $R_{\text{eff}}(X_i \leftrightarrow X_j) = \frac{2}{d} (1+o(1))$

Does this imply the Thm?

Or: If we know the $R(x,y)$ $\forall x \neq y$
 in a network, can we reconstruct
 the network? Is this stable?

Answer: Yes, up to loops

Lemma $k \geq 3$, $x > 0$ fixed, $\varepsilon > 0$ small

A network on $\{1, \dots, k\}$

If $|R(i,j) - x| < \varepsilon \quad \forall i \neq j \in [k]$

then

$$R(1, \{2, \dots, k\}) = \frac{kx}{2(k-1)} + O(\varepsilon)$$

Pf: Step 1: Knowing pairwise eff resistance
 we can recover all voltages, stable

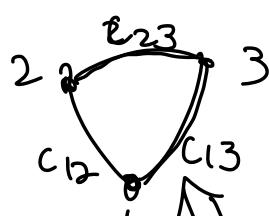
Indeed

$v_a(i,j) :=$ voltage at i when
 $a \rightarrow 0, j \rightarrow 1$

$$= p_i (\tau_j < \tau_a)$$

$$V_a(i,j) = \frac{R_{\text{eff}}(a \leftrightarrow j) + R_{\text{eff}}(a \leftrightarrow i) - \cancel{R_{\text{eff}}(i,j)}}{2 R_{\text{eff}}(a \leftrightarrow j)}$$

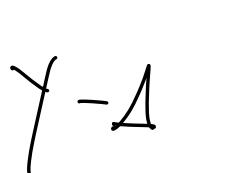
Sufficient to consider a network w/ 3 vertices



$$\underline{R_{\text{eff}}(1,2)} = \frac{1}{c_{12} + \frac{1}{\frac{1}{c_{13}} + \frac{1}{c_{23}}}}$$

$$=$$

$$=$$



$$\begin{aligned} R(1,2) &= r_1 + r_2 \\ R(2,3) &= r_2 + r_3 \\ R(1,3) &= r_1 + r_3 \end{aligned}$$

$$i=3$$

$$a, j = (1, 2)$$

$$\begin{pmatrix} 1 & 1 & 0 \\ 0 & 1 & 1 \\ 1 & 0 & 1 \end{pmatrix}$$

$$= r_3 = \frac{-R_{12} + R_{13} + R_{23}}{2}$$

Step 2: knowing the voltages
allows to recover the network.

$$\Delta(i,j) = \begin{cases} -C(i,j) & i \neq j \\ \pi(i) & i = j \end{cases}$$

$$a \in V \quad \Delta[a] = \begin{pmatrix} a \\ 4 \end{pmatrix}$$

$$\text{then } \underline{\Delta[a]} = \underline{\left(V_a(\cdot, \cdot) \right)^{-1}}$$

$$\underline{\text{Easy fact}} \quad \underline{\frac{1}{1+x}} = 1 + x + x^2 \dots$$

If A is invertible, E matrix with small entries
then $A+E$ is invertible and close coordinate
w/ $\|A^{-1}\|$

$\cdots x$ to A

$$V_a(\cdot, \cdot) \begin{pmatrix} x & & \\ & \ddots & x_{12} \\ x_{12} & \ddots & x \end{pmatrix}$$

$$(V_a)^{-1} = \begin{pmatrix} \frac{2(k-1)}{kx} & \dots & \text{(circle with } -\frac{2x}{k}) \\ \vdots & \ddots & \vdots \end{pmatrix}$$

$$R_{eff}(1, \{2, \dots, k\}) = \frac{1}{\sum_{j=2}^k C(1,j)} = \frac{1}{(k-1) \cdot \frac{2}{kx}}$$

